Exercises in System Level Programming (SLP) – Summer Term 2024

Exercise 6

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Interrupts



- Procedure of an interrupt (see 18-7):
 - o. Hardware sets required flag
 - 1. If interrupts are enabled and the interrupt is not masked, the interrupt controller interrupts the current execution
 - 2. Further interrupts are disabled
 - 3. Current program position is saved
 - 4. Address of the handler is read from the interrupt vector table and is then jumped to
 - 5. The interrupt handler is executed
 - 6. At the end of the interrupt handler, the instruction "return from interrupt" returns to the interrupted program and the re-enables of the interrupts

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Implementation of Interrupt Handlers



- For every interrupt, one bit for storing its state is available
- May lead to lost interrupts: An interrupt occurs during...
 - the execution of an interrupt handler (interrupts too fast)
 - disabled interrupts section (for synchronization of critical sections)
- This problem cannot be prevented in general
- → Risk minimization: Interrupt handler shall be as short as possible
 - Avoid any kind of loops and function calls
 - Do not use any blocking function (ADC/serial interface!)



- Timer
- Serial interface
- ADC (analog digital converter)
- External interrupts by level changes at certain I/O pins
 - Choice of level- or edge-triggered
 - Depend on the interrupt source
 - ⇒ ATmega328PB: 2 sources at the pins PD2 (INT0) and PD3 (INT1)
 - ⇒ BUTTON0 at PD2
 - ⇒ BUTTON1 at PD3
- More details in the ATmega328PB data sheet

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(Re-)Enabling Interrupts

- Interrupts can be enabled and disabled by special machine instructions
- The library avr-libc provides useful macros: #include <avr/interrupt.h>
 - sei() (set interrupt flag): enables interrupts (delayed by one instruction)
 - cli() (clear interrupt flag): disables all interrupts (immediately)
- Upon entering an interrupt handler, all interrupts are blocked automatically and unblocked again as soon as the handler is exited
- sei() should never be called from inside an interrupt handler
 - Potentially infinitely nested interrupt handlers
 - Possibility of a stack overflow
- \blacksquare At the start of the μ C, interrupts are disabled by default



- Interrupt sense control (ISC) bits of the ATmega328PB are located at the external interrupt control register A (EICRA)
- Position of the ISC-bits inside the register defined by macros

Interrupt INTO		Interrupt on	Interrupt INT1	
ISC01	ISC00	Interrupt on	ISC11	ISC10
0	0	low level	0	0
0	1	either edge	0	1
1	0	falling edge	1	0
1	1	rising edge	1	1

Example: Configuring INT1 of the ATmega328PB for a falling edge

```
/* the ISC-bits are located in the EICRA */
EICRA &= ~(1 << ISC10); // deleting ISC10
EICRA |= (1 << ISC11); // setting ISC11</pre>
```

(Un-)Masking Interrupts



- Single interrupts can be enabled (= unmasked) individually
 - ATmega328PB: External interrupt mask register (EIMSK)
- The bit positions inside of the register are defined by macros
 INTn
- A set bit enables the corresponding interrupt
- Example: Enabling the external interrupt INT1

```
O1 EIMSK |= (1 << INT1); // Unmask the external interrupt INT1
```



- Registering an interrupt handler is implemented by the C library
- Macro ISR (interrupt service routine) used for defining a handler function (#include <avr/interrupt.h>)
- Parameter: Desired vector
 - Available vectors: Refer to avr-libc documentation for avr/interrupt.h
 - Example: INT1_vect for external interrupt INT1
- Example: Implement handler for INT1

```
01 #include <avr/interrupt.h>
02
03 static volatile uint16_t counter = 0;
04
05 ISR(INT1_vect) {
    counter++;
07 }
```

Synchronization



- When an interrupt occurs, event = 1 is set
- Active waiting loop waits until event != 0
- Compiler detects that event is not changed within the loop
 - ⇒ the value of event is only loaded once from memory into a processor register
 - ⇒ endless loop

```
static uint8_t event = 0;
   ISR(INT0_vect) {
02
        event = 1;
03
04
05
   void main(void) {
06
        while(1) {
07
            while(event == 0) { /* wait for event */ }
08
            // handle event [...]
09
        }
10
11
```

Keyword volatile



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- When an interrupt occurs, event = 1 is set
- Active waiting loop waits until event != 0
- Compiler detects that event is not changed within the loop
 - ⇒ the value of event is only loaded once from memory into a processor register
 - \Rightarrow endless loop
- volatile enforces that the variable is loaded from memory before every access

```
static volatile uint8_t event = 0;
   ISR(INT0_vect) {
02
        event = 1;
03
04
05
   void main(void) {
06
        while(1) {
07
            while(event == 0) { /* wait for event */ }
80
            // handle event [...]
09
        }
10
```



- Missing volatile can lead to unexpected program execution
- Unnecessary use of volatile prevent certain compiler optimizations
- Correct use of volatile is task of the programmer!
- → Use volatile as rarely as possible but as often as required

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Lost Update



- Counting button presses that have to be processed
 - Incremented in the interrupt handler
 - Decremented in the main program to start the processing

```
static volatile uint8_t counter = 0;
01
02 ISR(INT0_vect) {
      counter++;
03
   }
04
05
   void main(void) {
06
      while(1) {
07
        if(counter > 0) {
08
09
10
          counter--;
11
          // handle pressed button
12
          // [...]
13
14
15
```



Interrupt handler

```
o5 ; C instruction: counter++
o6 lds r25, counter
o7 inc r25
o8 sts counter, r25
```

Line	counter	r24	r25
_	5		

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Lost Update



Main program

```
o1 ; C instruction: counter--;
o2 lds r24, counter
o3 dec r24
o4 sts counter, r24
```

```
o5 ; C instruction: counter++
o6 lds r25, counter
o7 inc r25
o8 sts counter, r25
```

Line	counter	r24	r25
_	5		
2	5	5	1



```
continuation counter coun
```

Interrupt handler

```
o5 ; C instruction: counter++
o6 lds r25, counter
o7 inc r25
o8 sts counter, r25
```

Line	counter	r24	r25
_	5		
2	5	5	_
3	5	4	_

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Lost Update



Main program

```
o1 ; C instruction: counter--;
02 lds r24, counter
03 dec r24
04 sts counter, r24
```

```
o5 ; C instruction: counter++
o6 lds r25, counter
o7 inc r25
o8 sts counter, r25
```

Line	counter	r24	r25
_	5		
2	5	5	-
3	5	4	_
6	5	4	5



Interrupt handler

```
05 ; C instruction: counter++
06 lds r25, counter
07 inc r25
08 sts counter, r25
```

Line	counter	r24	r25
_	5		
2	5	5	_
3	5	4	_
6	5	4	5
7	5	4	6

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Lost Update



Main program

```
o1 ; C instruction: counter--;
02 lds r24, counter
03 dec r24
04 sts counter, r24
```

```
o5 ; C instruction: counter++
o6 lds r25, counter
o7 inc r25
o8 sts counter, r25
```

Line	counter	r24	r25
_	5		
2	5	5	_
3	5	4	_
6	5	4	5
7	5	4	6
8	6	4	6



```
o1 ; C instruction: counter--;
o2 lds r24, counter
o3 dec r24
o4 sts counter, r24
```

Interrupt handler

```
o5 ; C instruction: counter++
o6 lds r25, counter
o7 inc r25
o8 sts counter, r25
```

Line	counter	r24	r25
_	5		
2	5	5	_
3	5	4	_
6	5	4	5
7	5	4	6
8	6	4	6
4	4	4	_

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16-Bit Access (Read Write)



- Concurrent use of 16 bit values (read write)
 - Incrementing in the interrupt handler
 - Reading in the main program

```
o1 static volatile uint16_t counter = 0;
   ISR(INTO_vect) {
02
      counter++;
03
   }
04
05
   void main(void) {
06
     if(counter > 300) {
07
        sb_led_on(YELLOW0);
80
      } else {
09
        sb_led_off(YELLOW0);
10
11
12
     // [...]
13
14
```



```
01  ; C instruction: if(counter>300)
02  lds  r22, counter
03  lds  r23, counter+1
04  cpi  r22, 0x2D
05  sbci  r23, 0x01
```

Interrupt handler

```
7; C instruction: counter++;
08 lds r24, counter
09 lds r25, counter+1
10 adiw r24,1
11 sts counter+1, r25
12 sts counter, r24
```

Line	counter	r22 & r23	r24 & r25
_	0x00ff	_	_

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16-Bit Access (Read Write)



Main program

```
01  ; C instruction: if(counter>300)
02  lds  r22, counter
03  lds  r23, counter+1
04  cpi  r22, 0x2D
05  sbci  r23, 0x01
```

```
7 ; C instruction: counter++;
8 lds r24, counter
9 lds r25, counter+1
10 adiw r24,1
11 sts counter+1, r25
12 sts counter, r24
```

Line	counter	r22 & r23	r24 & r25
_	0x00ff	_	_
2	0x00ff	0x??ff	_



```
01  ; C instruction: if(counter>300)
02  lds  r22, counter
03  lds  r23, counter+1
04  cpi  r22, 0x2D
05  sbci  r23, 0x01
```

Interrupt handler

```
7 ; C instruction: counter++;
08 lds r24, counter
09 lds r25, counter+1
10 adiw r24,1
11 sts counter+1, r25
12 sts counter, r24
```

Line	counter	r22 & r23	r24 & r25
_	0x00ff	_	_
2	0x00ff	0x??ff	_
8+9	0x00ff	0x??ff	0x00ff

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16-Bit Access (Read Write)



Main program

```
01 ; C instruction: if(counter>300)
02 lds r22, counter
03 lds r23, counter+1
04 cpi r22, 0x2D
05 sbci r23, 0x01
```

```
07 ; C instruction: counter++;
08 lds r24, counter
09 lds r25, counter+1
10 adiw r24,1
11 sts counter+1, r25
12 sts counter, r24
```

Line	counter	r22 & r23	r24 & r25
_	0x00ff	_	_
2	0x00ff	0x??ff	_
8+9	0x00ff	0x??ff	0x00ff
10	0x00ff	0x??ff	0x0100



```
01  ; C instruction: if(counter>300)
02  lds  r22, counter
03  lds  r23, counter+1
04  cpi  r22, 0x2D
05  sbci  r23, 0x01
```

Interrupt handler

```
07 ; C instruction: counter++;
08 lds r24, counter
09 lds r25, counter+1
10 adiw r24,1
11 sts counter+1, r25
12 sts counter, r24
```

Line	counter	r22 & r23	r24 & r25
_	0x00ff	_	_
2	0x00ff	0x??ff	_
8+9	0x00ff	0x??ff	0x00ff
10	0x00ff	0x??ff	0x0100
11+12	0x0100	0x??ff	0x0100

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16-Bit Access (Read Write)



Main program

```
01  ; C instruction: if(counter>300)
02  lds  r22, counter
03  lds  r23, counter+1
04  cpi  r22, 0x2D
05  sbci  r23, 0x01
```

```
07 ; C instruction: counter++;
08 lds r24, counter
09 lds r25, counter+1
10 adiw r24,1
11 sts counter+1, r25
12 sts counter, r24
```

Line	counter	r22 & r23	r24 & r25
_	0x00ff	_	_
2	0x00ff	0x??ff	_
8+9	0x00ff	0x??ff	0x00ff
10	0x00ff	0x??ff	0x0100
11+12	0x0100	0x??ff	0x0100
3	0x0100	0x01ff	_



```
continuation: if(counter>300)
c
```

Interrupt handler

```
7; C instruction: counter++;
8 lds r24, counter
9 lds r25, counter+1
10 adiw r24,1
11 sts counter+1, r25
12 sts counter, r24
```

Line	counter	r22 & r23	r24 & r25
_	0x00ff	_	_
2	0x00ff	0x??ff	_
8+9	0x00ff	0x??ff	0x00ff
10	0x00ff	0x??ff	0x0100
11+12	0x0100	0x??ff	0x0100
3	0x0100	0x01ff	_

 \Rightarrow In lines 4+5, the comparison uses 0x01ff (= 511) instead of 0x0100 (= 256). The comparison yields true and the LED is switched on.

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Blocking the Handling of Interrupt on the AVR



- Many more concurrency problems are possible
 - Non-atomic modification of shared data
 - Analysis of the problem by the application programmer
 - Choice of suitable synchronization primitives
- Solution here: Mutual exclusion by disabling interrupts
 - Blocking all interrupts: cli() and sei()
 - Disabling single interrupts (EIMSK-register)
- Problem: Interrupts can be lost during a blocked section
- ⇒ Critical sections have to be as short as possible



How can a lost update be prevented?

```
static volatile uint8_t counter = 0;
   ISR(INT0_vect) {
      counter++;
03
   }
04
05
   void main(void) {
06
     while(1) {
07
80
        if(counter > 0) {
09
          counter--;
10
11
          // handle pressed button
12
          // [...]
13
        }
14
15
16
```

16

Lost Update



How can a lost update be prevented?

```
static volatile uint8_t counter = 0;
   ISR(INTO_vect) {
02
      counter++;
03
   }
04
05
   void main(void) {
06
     while(1) {
07
        if(counter > 0) {
80
09
          cli();
          counter--;
10
          sei();
11
          // handle pressed button
12
          // [...]
13
        }
14
      }
15
16
```



How can a read-write anomaly be prevented?

```
static volatile uint16_t counter = 0;
    ISR(INT0_vect) {
02
03
      counter++;
   }
04
05
   void main(void) {
06
07
80
09
      if(counter > 300) {
10
11
        sb_led_on(YELLOW0);
12
      } else {
13
14
        sb_led_off(YELLOW0);
15
      }
16
17
      // [...]
18
19
```

16-Bit Access (Read Write)



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How can a read-write anomaly be prevented?

```
static volatile uint16_t counter = 0;
   ISR(INT0_vect) {
02
      counter++;
03
   }
04
05
   void main(void) {
06
      cli();
07
     uint16_t local_counter = counter;
80
      sei();
09
     if(local_counter > 300) {
10
11
        sb_led_on(YELLOW0);
12
      } else {
13
14
        sb_led_off(YELLOW0);
15
      }
16
17
      // [...]
18
19
```



How can a read-write anomaly be prevented?

```
static volatile uint16_t counter = 0;
01
    ISR(INT0_vect) {
02
03
      counter++;
   }
04
05
   void main(void) {
06
07
80
      cli();
09
      if(counter > 300) {
10
11
        sb_led_on(YELLOW0);
12
      } else {
13
14
        sb_led_off(YELLOW0);
15
16
      sei();
17
      // [...]
18
19
```

16-Bit Access (Read Write)



17

How can a read-write anomaly be prevented?

```
static volatile uint16_t counter = 0;
   ISR(INT0_vect) {
02
      counter++;
03
   }
04
05
   void main(void) {
06
07
80
      cli();
09
     if(counter > 300) {
10
        sei();
11
        sb_led_on(YELLOW0);
12
      } else {
13
        sei();
14
        sb_led_off(YELLOW0);
15
      }
16
17
     // [...]
18
19
```

Power-Saving Modes

Power-Saving Modes of AVR Processors



- AVR-based devices are often powered by batteries (e.g. remotes)
- Saving energy can drastically extend the life span
- AVR processors support multiple power-saving modes
 - Deactivating functional units
 - Different "depths" of sleep
 - Only active functional units can wake up the CPU
- Default mode: Idle
 - CPU clock is stopped
 - Nor more memory accesses
 - Hardware (timer, external interrupts, ADC, etc.) are still active
- Documentation in ATmega328PB data sheet



- Support from the avr-libc: (#include <avr/sleep.h>)
 - sleep_enable() enables the sleep mode
 - sleep_cpu() enters the sleep mode
 - sleep_disable() disables the sleep mode
 - set_sleep_mode(uint8_t mode) configures the used mode
- Documentation of avr/sleep.h in avr-libc documentation

```
#include <avr/sleep.h>

set_sleep_mode(SLEEP_MODE_IDLE); // use idle mode

sleep_enable(); // activate sleep mode

sleep_cpu(); // enter sleep mode

sleep_disable(); // recommended: deactivate sleep mode

→ afterwards
```

Lost Wakeup



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- Sleeping beauty (german: Dornröschenschlaf)
 - ⇒ **Problem:** There is exactly one interrupt

Main program

```
O1 ISR(TIMER1_COMPA_vect) {
O2    event = 1;
O3 }
```



- Sleeping beauty (german: Dornröschenschlaf)
 - ⇒ **Problem:** There is exactly one interrupt

```
sleep_enable();
01
    event = 0;
02
03
04
   while(!event) {
05
        ∮ Interrupt ∮
06
        sleep_cpu();
07
08
09
10
11
   sleep_disable();
```

Interrupt handler

```
O1   ISR(TIMER1_COMPA_vect) {
        event = 1;
        o3   }
```

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Lost Wakeup



- Sleeping beauty (german: Dornröschenschlaf)
 - ⇒ **Problem:** There is exactly one interrupt
 - ⇒ **Solution:** Disable interrupts during the critical area

Main program

```
sleep_enable();
    event = 0;
02
03
    cli();
04
   while(!event) {
05
        sei();
06
        sleep_cpu();
07
        cli();
80
09
    sei();
10
11
    sleep_disable();
```

```
O1 ISR(TIMER1_COMPA_vect) {
O2     event = 1;
O3 }
```



- Sleeping beauty (german: Dornröschenschlaf)
 - ⇒ **Problem:** There is exactly one interrupt
 - ⇒ **Solution:** Disable interrupts during the critical area

```
sleep_enable();
01
    event = 0;
02
03
   cli();
04
   while(!event) {
05
        sei(); ∮ Interrupt ∮
06
        sleep_cpu();
07
        cli();
08
09
   sei();
10
11
   sleep_disable();
```

Interrupt handler

```
o1 ISR(TIMER1_COMPA_vect) {
     event = 1;
     o3 }
```

⇒ What if the interrupt occurs between lines 6 and 7?

Lost Wakeup



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- Sleeping beauty (german: Dornröschenschlaf)
 - ⇒ **Problem:** There is exactly one interrupt
 - ⇒ **Solution:** Disable interrupts during the critical area

Main program

```
sleep_enable();
   event = 0;
02
03
   cli();
04
   while(!event) {
05
        sei(); ∮ Interrupt ∮
06
        sleep_cpu();
07
        cli();
08
09
   sei();
10
11
   sleep_disable();
```

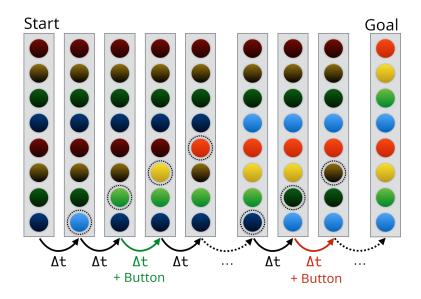
```
O1 ISR(TIMER1_COMPA_vect) {
O2     event = 1;
O3 }
```

Assignment: Dexterity Game

Assignment: Dexterity Game (1)



- Game cursor moves over the LED strip and inverts (toggles) the state of the LED
- LED state is retained if the button is pressed
- Goal: Switch on all LEDs





After each level, a winning sequence is displayed via the LEDs

```
void main(void) {
      // Initialisation
02
     // [...]
03
04
     while(1) {
05
        // starting level
06
        // [...]
07
80
        // show win sequence
09
        // [...]
10
11
        // update level
12
        // [...]
13
14
15
```

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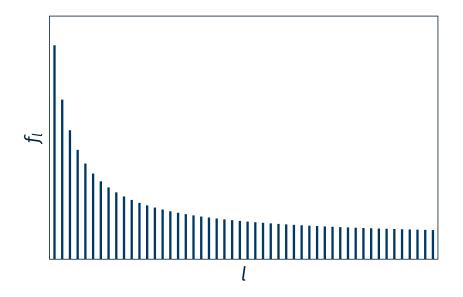
Detect a Button Press



- Goals:
 - Edge detection in hardware
 - Handle events using interrupts
 - No use of the libspicboard
- Details:
 - BUTTON0 is wired to PD2
 - Configure PD2 as input (with activated pull-up resistor)
 - PD2 is input of INT0
 - Which level/edge has to be configured for the interrupt?
 - How does a minimal interrupt handler for this assignment look like?



- Speed of the game determines its difficulty
 - ⇒ Passive waiting with the timer module of the libspicboard
- Difficulty increases with each level l
- Speed converges to a maximum
 - \Rightarrow Series of waiting times: $f_l = \frac{a}{l} + b$ (a and b are constants)



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Hands-on: Simple Interrupt Counter



- Counting activations of BUTTON0 (PD2)
- Detect activation with the help of interrupts
- Output the current counter value using the 7-segment display
- Enter a CPU sleeping state whenever the value is even
- "Standby" LED switched on during the sleep mode (BLUE0)
- Hints:
 - Detection of the activation without the libspicboard
 - PD2/BUTTON0 is the input of INT0
 - Interrupt on a falling edge:
 - EICRA(ISC00) = 0
 - EICRA(ISC01) = 1
 - 7-segment display needs regular interrupts to display values