Exercises in System Level Programming (SLP) – Summer Term 2024

Exercise 6

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Interrupts

Interrupts



- Procedure of an interrupt (see 18-7):
 - o. Hardware sets required flag
 - 1. If interrupts are enabled and the interrupt is not masked, the interrupt controller interrupts the current execution
 - 2. Further interrupts are disabled
 - 3. Current program position is saved
 - 4. Address of the handler is read from the interrupt vector table and is then jumped to
 - 5. The interrupt handler is executed
 - 6. At the end of the interrupt handler, the instruction "return from interrupt" returns to the interrupted program and the re-enables of the interrupts

Implementation of Interrupt Handlers



- For every interrupt, one bit for storing its state is available
- May lead to lost interrupts: An interrupt occurs during...
 - the execution of an interrupt handler (interrupts too fast)
 - disabled interrupts section (for synchronization of critical sections)
- This problem cannot be prevented in general
- → Risk minimization: Interrupt handler shall be as short as possible
 - Avoid any kind of loops and function calls
 - Do not use any blocking function (ADC/serial interface!)

Interrupts on the AVR



- Timer
- Serial interface
- ADC (analog digital converter)
- External interrupts by level changes at certain I/O pins
 - Choice of level- or edge-triggered
 - Depend on the interrupt source
 - ⇒ ATmega328PB: 2 sources at the pins PD2 (INT0) and PD3 (INT1)
 - ⇒ BUTTON0 at PD2
 - ⇒ BUTTON1 at PD3
- More details in the ATmega328PB data sheet

(Re-)Enabling Interrupts



- Interrupts can be enabled and disabled by special machine instructions
- The library avr-libc provides useful macros: #include <avr/interrupt.h>
 - sei() (set interrupt flag): enables interrupts (delayed by one instruction)
 - cli() (clear interrupt flag): disables all interrupts (immediately)
- Upon entering an interrupt handler, all interrupts are blocked automatically and unblocked again as soon as the handler is exited
- sei() should never be called from inside an interrupt handler
 - Potentially infinitely nested interrupt handlers
 - Possibility of a stack overflow
- \blacksquare At the start of the μ C, interrupts are disabled by default

Configuring Interrupts



- Interrupt sense control (ISC) bits of the ATmega328PB are located at the external interrupt control register A (EICRA)
- Position of the ISC-bits inside the register defined by macros

Interrupt INT0		Interrupt on	Interrupt INT1	
_ISC01	ISC00	Interrupt on	ISC11	ISC10
0	0	low level	0	0
0	1	either edge	0	1
1	0	falling edge	1	0
1	1	rising edge	1	1

Example: Configuring INT1 of the ATmega328PB for a falling edge

```
01  /* the ISC-bits are located in the EICRA */
02  EICRA δ= ~(1 << ISC10); // deleting ISC10
03  EICRA |= (1 << ISC11); // setting ISC11</pre>
```

(Un-)Masking Interrupts



- Single interrupts can be enabled (= unmasked) individually
 - ATmega328PB: External interrupt mask register (EIMSK)
- The bit positions inside of the register are defined by macros INTn
- A set bit enables the corresponding interrupt
- Example: Enabling the external interrupt INT1

O1 EIMSK |= (1 << INT1); // Unmask the external interrupt INT1



- Registering an interrupt handler is implemented by the C library
- Macro ISR (interrupt service routine) used for defining a handler function (#include <avr/interrupt.h>)
- Parameter: Desired vector
 - Available vectors: Refer to avr-libc documentation for avr/interrupt.h
 - Example: INT1_vect for external interrupt INT1
- Example: Implement handler for INT1

```
01 #include <avr/interrupt.h>
02
03 static volatile uint16_t counter = 0;
04
05 ISR(INT1_vect) {
06    counter++;
07 }
```

Synchronization

Keyword volatile



- When an interrupt occurs, event = 1 is set
- Active waiting loop waits until event != 0
- Compiler detects that event is not changed within the loop
 - ⇒ the value of event is only loaded once from memory into a processor register
 - ⇒ endless loop

```
o1 static uint8_t event = 0;
   ISR(INTO_vect) {
03
       event = 1;
04
05
   void main(void) {
06
       while(1) {
07
            while(event == 0) { /* wait for event */ }
08
            // handle event [...]
09
10
11
```

Keyword volatile



- When an interrupt occurs, event = 1 is set
- Active waiting loop waits until event != 0
- Compiler detects that event is not changed within the loop
 - ⇒ the value of event is only loaded once from memory into a processor register
 - \Rightarrow endless loop
- volatile enforces that the variable is loaded from memory before every access

Usage of volatile



- Missing volatile can lead to unexpected program execution
- Unnecessary use of volatile prevent certain compiler optimizations
- Correct use of volatile is task of the programmer!
- → Use volatile as rarely as possible but as often as required



- Counting button presses that have to be processed
 - Incremented in the interrupt handler
 - Decremented in the main program to start the processing

```
static volatile uint8 t counter = 0;
   ISR(INT0 vect) {
02
03
     counter++;
04
05
   void main(void) {
     while(1) {
07
08
        if(counter > 0) {
09
          counter--;
10
11
          // handle pressed button
12
         // [...]
13
14
15
16
```



Main program

```
O1 ; C instruction: counter--;
O2 lds r24, counter
O3 dec r24
O4 sts counter, r24
```

```
o5 ; C instruction: counter++
o6 lds r25, counter
o7 inc r25
o8 sts counter, r25
```

Line	counter	r24	r25
_	5		



Main program

```
o1 ; C instruction: counter--;
o2 lds r24, counter
o3 dec r24
o4 sts counter, r24
```

```
o5 ; C instruction: counter++
o6 lds r25, counter
o7 inc r25
o8 sts counter, r25
```

Line	counter	r24	r25
_	5		
2	5	5	_



Main program

```
01 ; C instruction: counter--;
02 lds r24, counter
03 dec r24
04 sts counter, r24
```

```
o5 ; C instruction: counter++
o6 lds r25, counter
o7 inc r25
o8 sts counter, r25
```

Line	counter	r24	r25
_	5		
2	5	5	_
3	5	4	_



Main program

```
o1 ; C instruction: counter--;
o2 lds r24, counter
o3 dec r24
o4 sts counter, r24
```

```
o5 ; C instruction: counter++
o6 lds r25, counter
o7 inc r25
o8 sts counter, r25
```

Line	counter	r24	r25
_	5		
2	5	5	_
3	5	4	_
6	5	4	5



Main program

```
01 ; C instruction: counter--;
02 lds r24, counter
03 dec r24
04 sts counter, r24
```

```
o5 ; C instruction: counter++
o6 lds r25, counter
o7 inc r25
o8 sts counter, r25
```

Line	counter	r24	r25
_	5		
2	5	5	_
3	5	4	_
6	5	4	5
7	5	4	6



Main program

```
O1 ; C instruction: counter--;
O2 lds r24, counter
O3 dec r24
O4 sts counter, r24
```

```
o5 ; C instruction: counter++
o6 lds r25, counter
o7 inc r25
o8 sts counter, r25
```

Line	counter	r24	r25
-	5		
2	5	5	_
3	5	4	_
6	5	4	5
7	5	4	6
8	6	4	6



Main program

```
01 ; C instruction: counter--;
02 lds r24, counter
03 dec r24
04 sts counter, r24
```

```
o5 ; C instruction: counter++
o6 lds r25, counter
o7 inc r25
o8 sts counter, r25
```

Line	counter	r24	r25
-	5		
2	5	5	_
3	5	4	_
6	5	4	5
7	5	4	6
8	6	4	6
4	4	4	_



- Concurrent use of 16 bit values (read write)
 - Incrementing in the interrupt handler
 - Reading in the main program

```
static volatile uint16_t counter = 0;
   ISR(INT0_vect) {
     counter++;
03
04
05
   void main(void) {
     if(counter > 300) {
07
        sb led on(YELLOW0);
08
09
     } else {
        sb_led_off(YELLOW0);
10
11
12
     // [...]
13
14
```



Main program

```
01 ; C instruction: if(counter>300)
02 lds r22, counter
03 lds r23, counter+1
04 cpi r22, 0x2D
05 sbci r23, 0x01
```

```
o7 ; C instruction: counter++;
08 lds r24, counter
09 lds r25, counter+1
10 adiw r24,1
11 sts counter+1, r25
12 sts counter, r24
```

Line	counter	r22 & r23	r24 & r25
_	0x00ff	_	_



Main program

```
01 ; C instruction: if(counter>300)
02 lds r22, counter
03 lds r23, counter+1
04 cpi r22, 0x2D
05 sbci r23, 0x01
```

```
07 ; C instruction: counter++;
08 lds r24, counter
09 lds r25, counter+1
10 adiw r24,1
11 sts counter+1, r25
12 sts counter, r24
```

Line	counter	r22 & r23	r24 & r25
_	0x00ff	_	_
2	0x00ff	0x??ff	_



Main program

```
01 ; C instruction: if(counter>300)
02 lds r22, counter
03 lds r23, counter+1
04 cpi r22, 0x2D
05 sbci r23, 0x01
```

```
07 ; C instruction: counter++;
08 lds r24, counter
09 lds r25, counter+1
10 adiw r24,1
11 sts counter+1, r25
12 sts counter, r24
```

Line	counter	r22 & r23	r24 & r25
_	0x00ff	_	_
2	0x00ff	0x??ff	_
8+9	0x00ff	0x??ff	0x00ff



Main program

```
O1 ; C instruction: if(counter>300)
O2 lds r22, counter
O3 lds r23, counter+1
O4 cpi r22, 0x2D
O5 sbci r23, 0x01
```

```
o7 ; C instruction: counter++;
08 lds r24, counter
09 lds r25, counter+1
10 adiw r24,1
11 sts counter+1, r25
12 sts counter, r24
```

Line	counter	r22 & r23	r24 & r25
_	0x00ff	_	_
2	0x00ff	0x??ff	_
8+9	0x00ff	0x??ff	0x00ff
10	0x00ff	0x??ff	0x0100



Main program

```
01 ; C instruction: if(counter>300)
02 lds r22, counter
03 lds r23, counter+1
04 cpi r22, 0x2D
05 sbci r23, 0x01
```

```
07 ; C instruction: counter++;
08 lds r24, counter
09 lds r25, counter+1
10 adiw r24,1
11 sts counter+1, r25
12 sts counter, r24
```

Line	counter	r22 & r23	r24 & r25
_	0x00ff	_	_
2	0x00ff	0x??ff	_
8+9	0x00ff	0x??ff	0x00ff
10	0x00ff	0x??ff	0x0100
11+12	0x0100	0x??ff	0x0100



Main program

```
01 ; C instruction: if(counter>300)
02 lds r22, counter
03 lds r23, counter+1
04 cpi r22, 0x2D
05 sbci r23, 0x01
```

```
07 ; C instruction: counter++;
08 lds r24, counter
09 lds r25, counter+1
10 adiw r24,1
11 sts counter+1, r25
12 sts counter, r24
```

Line	counter	r22 & r23	r24 & r25
_	0x00ff	_	_
2	0x00ff	0x??ff	_
8+9	0x00ff	0x??ff	0x00ff
10	0x00ff	0x??ff	0x0100
11+12	0x0100	0x??ff	0x0100
3	0x0100	0x01ff	_



Main program

```
01 ; C instruction: if(counter>300)
02 lds r22, counter
03 lds r23, counter+1
04 cpi r22, 0x2D
05 sbci r23, 0x01
```

Interrupt handler

```
07 ; C instruction: counter++;
08 lds r24, counter
09 lds r25, counter+1
10 adiw r24,1
11 sts counter+1, r25
12 sts counter, r24
```

Line	counter	r22 & r23	r24 & r25
_	0x00ff	_	_
2	0x00ff	0x??ff	_
8+9	0x00ff	0x??ff	0x00ff
10	0x00ff	0x??ff	0x0100
11+12	0x0100	0x??ff	0x0100
3	0x0100	0x01ff	_

⇒ In lines 4+5, the comparison uses 0x01ff (= 511) instead of 0x0100 (= 256). The comparison yields true and the LED is switched on.

Blocking the Handling of Interrupt on the AVR



- Many more concurrency problems are possible
 - Non-atomic modification of shared data
 - Analysis of the problem by the application programmer
 - Choice of suitable synchronization primitives
- Solution here: Mutual exclusion by disabling interrupts
 - Blocking all interrupts: cli() and sei()
 - Disabling single interrupts (EIMSK-register)
- Problem: Interrupts can be lost during a blocked section
- ⇒ Critical sections have to be as short as possible



How can a lost update be prevented?

```
static volatile uint8 t counter = 0;
   ISR(INT0 vect) {
02
     counter++;
03
04
05
06
   void main(void) {
     while(1) {
07
       if(counter > 0) {
08
09
          counter--;
10
11
12
          // handle pressed button
          // [...]
13
14
15
16
```



How can a lost update be prevented?

```
static volatile uint8 t counter = 0;
   ISR(INT0 vect) {
02
     counter++;
03
04
05
   void main(void) {
     while(1) {
07
       if(counter > 0) {
08
09
          cli();
          counter--;
10
          sei();
11
          // handle pressed button
12
          // [...]
13
14
15
16
```



```
static volatile uint16_t counter = 0;
   ISR(INT0_vect) {
03
     counter++;
04
05
   void main(void) {
06
07
08
09
     if(counter > 300) {
10
11
        sb_led_on(YELLOW0);
12
     } else {
13
14
15
        sb_led_off(YELLOW0);
16
17
     // [...]
18
19
```



```
static volatile uint16_t counter = 0;
   ISR(INT0_vect) {
03
     counter++;
04
05
   void main(void) {
06
     cli();
07
     uint16 t local counter = counter;
08
     sei();
09
     if(local counter > 300) {
10
11
        sb_led_on(YELLOW0);
12
     } else {
13
14
        sb_led_off(YELLOW0);
15
16
17
     // [...]
18
19
```



```
static volatile uint16_t counter = 0;
   ISR(INT0_vect) {
03
      counter++;
04
05
   void main(void) {
06
07
08
     cli();
09
      if(counter > 300) {
10
11
        sb_led_on(YELLOW0);
12
      } else {
13
14
        sb_led_off(YELLOW0);
15
16
      sei();
17
      // [...]
18
19
```



```
static volatile uint16_t counter = 0;
   ISR(INT0_vect) {
03
      counter++;
04
05
   void main(void) {
06
07
08
      cli();
09
      if(counter > 300) {
10
        sei();
11
        sb_led_on(YELLOW0);
12
      } else {
13
        sei();
14
15
        sb_led_off(YELLOW0);
16
17
     // [...]
18
19
```

Power-Saving Modes

Power-Saving Modes of AVR Processors



- AVR-based devices are often powered by batteries (e.g. remotes)
- Saving energy can drastically extend the life span
- AVR processors support multiple power-saving modes
 - Deactivating functional units
 - Different "depths" of sleep
 - Only active functional units can wake up the CPU
- Default mode: Idle
 - CPU clock is stopped
 - Nor more memory accesses
 - Hardware (timer, external interrupts, ADC, etc.) are still active
- Documentation in ATmega328PB data sheet

Usage of the Sleep Modes



- Support from the avr-libc: (#include <avr/sleep.h>)
 - sleep_enable() enables the sleep mode
 - sleep_cpu() enters the sleep mode
 - sleep_disable() disables the sleep mode
 - set_sleep_mode(uint8_t mode) configures the used mode
- Documentation of avr/sleep.h in avr-libc documentation

```
#include <avr/sleep.h>

set_sleep_mode(SLEEP_MODE_IDLE); // use idle mode

sleep_enable(); // activate sleep mode

sleep_cpu(); // enter sleep mode

sleep_disable(); // recommended: deactivate sleep mode

afterwards
```



- Sleeping beauty (german: Dornröschenschlaf)
 - ⇒ **Problem:** There is exactly one interrupt

Main program

```
sleep enable();
   event = 0;
02
03
04
   while(!event) {
05
06
        sleep_cpu();
07
08
09
10
11
   sleep disable();
```

Interrupt handler



- Sleeping beauty (german: Dornröschenschlaf)
 - ⇒ **Problem:** There is exactly one interrupt

Main program

```
sleep_enable();
   event = 0;
02
03
04
    while(!event) {
05

    Interrupt ∮

06
        sleep_cpu();
07
08
09
10
11
   sleep disable();
```

Interrupt handler

```
01 ISR(TIMER1_COMPA_vect) {
02     event = 1;
03 }
```



- Sleeping beauty (german: Dornröschenschlaf)
 - ⇒ **Problem:** There is exactly one interrupt
 - ⇒ **Solution:** Disable interrupts during the critical area

Main program

```
sleep enable();
   event = 0;
02
03
  cli();
   while(!event) {
        sei();
06
       sleep_cpu();
07
08
        cli();
09
   sei();
11
   sleep disable();
```

Interrupt handler

```
o1 ISR(TIMER1_COMPA_vect) {
    event = 1;
    o3 }
```



- Sleeping beauty (german: Dornröschenschlaf)
 - ⇒ **Problem:** There is exactly one interrupt
 - ⇒ **Solution:** Disable interrupts during the critical area

Main program

```
sleep enable();
   event = 0;
02
03
04 cli();
   while(!event) {
        sei(); / Interrupt /
06
       sleep_cpu();
07
       cli();
08
09
   sei();
11
   sleep disable();
```

Interrupt handler

```
o1 ISR(TIMER1_COMPA_vect) {
    event = 1;
    o3 }
```

⇒ What if the interrupt occurs between lines 6 and 7?



- Sleeping beauty (german: Dornröschenschlaf)
 - ⇒ **Problem:** There is exactly one interrupt
 - ⇒ **Solution:** Disable interrupts during the critical area

Main program

```
sleep enable();
   event = 0;
02
03
  cli();
   while(!event) {
       sei(); # Interrupt #
06
       sleep_cpu();
07
       cli();
08
09
   sei();
11
   sleep disable();
```

Interrupt handler

```
O1 ISR(TIMER1_COMPA_vect) {
O2    event = 1;
O3 }
```

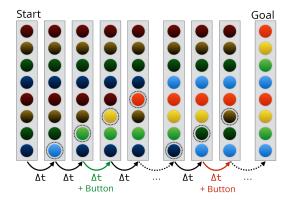
⇒ **Solution:** sei() is executed atomically with next line

Assignment: Dexterity Game

Assignment: Dexterity Game (1)



- Game cursor moves over the LED strip and inverts (toggles) the state of the LED
- LED state is retained if the button is pressed
- Goal: Switch on all LEDs



Assignment: Dexterity Game (2)



After each level, a winning sequence is displayed via the LEDs

```
void main(void) {
02
     // Initialisation
     // [...]
03
04
     while(1) {
05
       // starting level
06
07
       // [...]
08
       // show win sequence
09
       // [...]
10
11
       // update level
12
13
       // [...]
14
15
```

Detect a Button Press



Goals:

- Edge detection in hardware
- Handle events using interrupts
- No use of the libspicboard

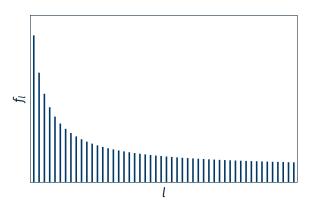
Details:

- BUTTON0 is wired to PD2
- Configure PD2 as input (with activated pull-up resistor)
- PD2 is input of INT0
- Which level/edge has to be configured for the interrupt?
- How does a minimal interrupt handler for this assignment look like?

Difficulty



- Speed of the game determines its difficulty
 - ⇒ Passive waiting with the timer module of the libspicboard
- Difficulty increases with each level *l*
- Speed converges to a maximum
 - \Rightarrow Series of waiting times: $f_l = \frac{a}{l} + b$ (a and b are constants)



Hands-on: Simple Interrupt Counter



- Counting activations of BUTTON0 (PD2)
- Detect activation with the help of interrupts
- Output the current counter value using the 7-segment display
- Enter a CPU sleeping state whenever the value is even
- "Standby" LED switched on during the sleep mode (BLUE0)
- Hints:
 - Detection of the activation without the libspicboard
 - PD2/BUTTON0 is the input of INT0
 - Interrupt on a falling edge:
 - EICRA(ISC00) = 0
 - EICRA(ISC01) = 1
 - 7-segment display needs regular interrupts to display values