Exercises in System Level Programming (SLP) – Summer Term 2024

Exercise 7

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Presentation Assignment 3

AVR Timer

Timer: Motivation



- Common task for *µ*Controller programming:
 - Regularly updating an output (e.g. frame rate)
 - Regularly reading of a value (e.g. serial console)
 - Pulse width modulation (PWM)
 - Passive waiting

• ...

 \Rightarrow Implementation using a *timer*

- A timer modifies a counter in every cycle
 - Increment (default)
 - Decrement
- When a previously configured event occurs, an interrupt is generated
 - Counter reaches a specific value
 - Counter overflows
 - (external event occurs)
- The ATmega328PB provides 5 different timers:
 - TIMER{0,2}: 8-bit counter
 - TIMER{1,3,4}: 16-bit counter
- \Rightarrow For all exercise tasks: TIMER0
- \Rightarrow Used by the libspicboard: TIMER{1,2,4}

Timer: Configuration (Timer clock speed)

How fast does the timer run:

- TCCR0B: TC0 control register B
- CSxx: Clock select bits
- Prescaler: Amount of CPU cycles until the counter is incremented
- What happens when the CPU enters a sleeping state?

CS02	CS01	CS00	Description	
0	0	0	Timer off	
0	0	1	prescaler 1	
0	1	0	prescaler 8	
0	1	1	prescaler 64	
1	0	0	prescaler 256	
1	0	1	prescaler 1024 Ext. clock (falling edge) Ext. clock (rising edge)	
1	1	0		
1	1	1		



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Timer: Configuration (Trigger Event)

• When does the timer trigger an interrupt:

- Overflow: When the counter flows over
- Match: When the counter reaches a specific value
 - ⇒ Register OCR0A (TIMER0 Output Compare Register A)
 - \Rightarrow Register OCR0B (TIMER0 Output Compare Register B)
- Interrupts can be unmasked individually
- TIMSK0: TIMER0 Interrupt Mask Register

Bit	ISR	Description	
TOIE0	TIMER0_OVF_vect	TIMER0 Overflow (Interrupt Enable)	
OCIE0A TIMER0_COMPA_vect		TIMER0 Output Compare A ()	
OCIE0B	TIMER0_COMPB_vect	TIMER0 Output Compare B ()	

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01	<pre>ISR(TIMER0_OVF_vect) {</pre>
02	// []
03	}
04	
05	<pre>static void init(void) {</pre>
06	// Activate overflow interrupt
07	TIMSK0 = (1 << TOIE0);
08	
09	// []
10	}

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■ Reminder: prescaler ∈ {1, 8, 64, 256, 1024}

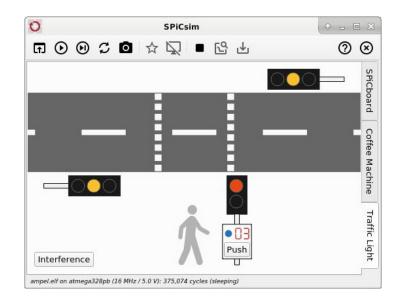
• Example:

- 8-bit timer with overflow interrupt
- CPU frequency: 16 MHz (ATmega328PB)
- Goal: Count with a cycle of length 1 $\rm s$
- \Rightarrow Which prescaler is the most resource efficient?
- \Rightarrow How many overflow interrupts are required until 1s has passed?
- \Rightarrow How big is the error that we have to accept?

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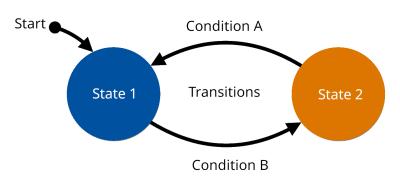
Task: Traffic Light





 Implementation of a (pedestrian) traffic light with waiting-time display

Finite State Machines

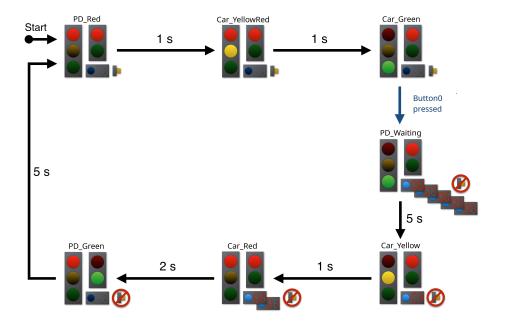


- States with specific attributes; well-defined initial state
- Transition depends on certain conditions

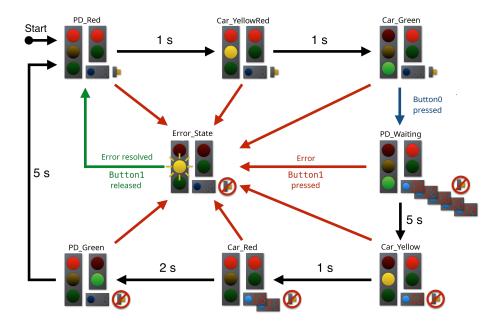
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Traffic Light as a Finite State Machine





Traffic Light as a Finite State Machine



- Using states with hardcoded integer values is prone to errors
 - Hard to memorize
 - Range of value cannot easily be restricted
- Better enum:

02

```
01 enum state { STATE_RED, STATE_YELLOW, STATE_GREEN };
02
```

```
o3 enum state my_state = STATE_RED;
```

• With typedef even more readable:

```
01 typedef enum { STATE_RED, STATE_YELLOW, STATE_GREEN } state;
```

o3 state my_state = STATE_RED;

Choosing States: switch-case Instruction

01	<pre>switch (my_state) {</pre>
02	case STATE_RED:
03	
04	break;
05	case STATE_YELLOW:
06	
07	break;
08	case STATE_GREEN:
09	
10	break;
11	default:
12	// maybe invalid state
13	
14	}

- Avoid any if-else-cascades
- switch-expression has to be an integer (or even better: enum)
- Do not forget the break-instruction!
- Ideal for handling systems with different states
 - \Rightarrow Implementation of finite state machines



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- Each transition is triggered by an interrupt
 - Configure BUTTON0 and BUTTON1 as interrupt inputs
 - \Rightarrow Which edge should trigger the interrupt?
 - Configure TIMER0 (interval: 1 second)
- Do not use the timer module of the libspicboard when submitting
 - \Rightarrow However, its use can be helpful for debugging

Hints

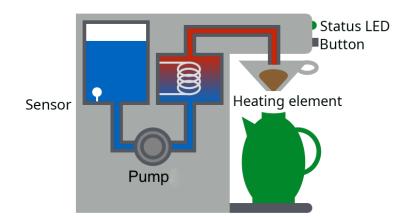
Hints:

- Implement each function exactly as specified in the description (reference implementation available)
- Model presses of the buttons and alarms as events
- Wait passively for all interrupts
- "Deactivate" the button by simply ignoring its interrupt (It is not necessary to modify the interrupt configuration)
- Mapping to a finite state machine can be useful
- Usage of volatile always needs a reason

Hands-on: Coffee Machine

Screencast: https://www.video.uni-erlangen.de/clip/id/17647

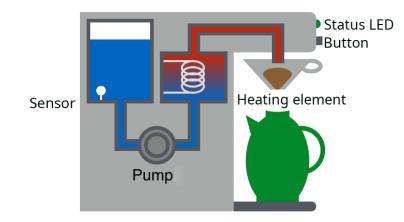
Hands-on: Coffee Machine (1)



- Learning goals:
 - Finite state machines
 - Timers and alarms
 - Interrupts & sleep modes

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Wiring:

- Pump & heating: Port D, Pin 5 (active-low)
- Button: INT0 an Port D, Pin 2 (active-low)
- Sensor: INT1 an Port D, Pin 3 (water: high; no water: low)
- State LED:
 - BLUE0: STANDBY
 - GREEN0: ACTIVE
 - RED0: NO_WATER

Hands-on: Coffee Machine (2)

STANDBY

- Machine is switched off
- Pump and heating are off
- User can start making coffee by pressing the button
- Initial state

ACTIVE

- Machine is switched on
- Pump and heating are on
- Water tank is not empty
- User can stop the machine by pressing the button

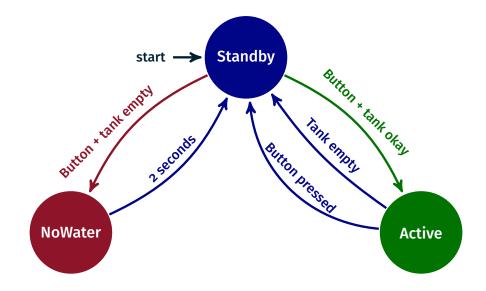
NO_WATER

- Coffee machine shows that not enough water is in the tank
- Pump and heating are off
- Time period: 2 seconds

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Hands-on: Coffee Machine (2)





- Hints:
 - Pressed button & change of water level by interrupts
 - State LED: void setLEDState(state_t state)
 - Waiting phases can be implemented using the single-shot alarms
 - During waiting phases always enter a power saving mode

Hands-on: Coffee Machine (3)

DDRx Configuration of pin i of port x as in-/output

- Bit i = $1 \rightarrow Pinias$ output
- Bit i = $0 \rightarrow Pinias input$

PORTx Mode of operation **depends on DDRx**:

- If pin i is configured as output, then bit i in the PORTx register controls whether a high level or a low level has to be generated at pin i
 - Bit i = $1 \rightarrow \text{high level at pin i}$
 - Bit i = $0 \rightarrow low level a pin i$
- If pin i is configured as input, then the internal pull-up resistor can be activated
 - Bit i = 1 → pull-up resistor at pin i (level is pulled high)
 - Bit i = $0 \rightarrow pin i configured as tri-state$

PINx Bit i returns the current level of pin i at port x (read only)

AP

- Interrupt sense control (ISC) bits of the ATmega328PB are located at the external interrupt control register A (EICRA)
- Position of the ISC-bits inside the register defined by macros

Interrupt INT0		Interrupt on	Interrupt INT1	
ISC01	ISC00	Interrupt on	ISC11	ISC10
0	0	low level	0	0
0	1	either edge	0	1
1	0	falling edge	1	0
1	1	rising edge	1	1

- ATmega328PB: External interrupt mask register (EIMSK)
- The position of the bits in this register is also defined by macros INTn