

# Verlässliche Echtzeitsysteme

## Übungen zur Vorlesung

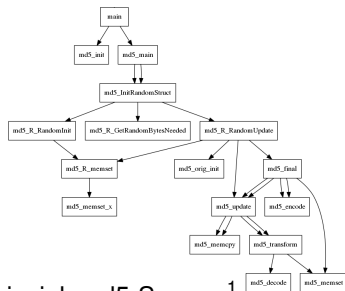
### Statische Stackbedarfsanalyse

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```

1  /* Objective function */
2  max: +16 md5_orig_init +64 md5_update \
3      +64 md5_final +16 md5_memset \
4      +208 md5_transform +16 md5_encode ...;
5
6
7  /* Constraints */
8  +main = 1;
9  +md5_init +md5_main <= +main;
10 ...
  
```

■ Beispiel: md5-Summe<sup>1</sup>

■ Vorgehen

1. Callgraph bestimmen
2. Stackbedarf einzelner Funktionen (gcc -fstack-usage)
3. ILP<sup>2</sup> aufstellen (Nebenbedingungen aus 1., Kosten aus 2. verwenden)
4. ILP z.B. mittels lp\_solve  $\rightsquigarrow$  **maximaler Stackbedarf**

<sup>1</sup><https://github.com/tacle/tacle-bench/>

<sup>2</sup>Integer Linear Program (dt. ganzzahliges lineares Programm)

## Optimierungsziel

- Jeder Stapelrahmen einer Funktion  $f$  hat eine Größe  $size$
- Jede Funktion kann auf einem Pfad ein- oder mehrfach (Rekursion), insgesamt  $n$ -fach auf dem Stapel vorkommen
- Gesucht: Fluss durch den Aufrufgraphen, welcher Stapelbedarf maximiert
- Dabei müssen **Flussbedingungen** eingehalten werden
  - Aufruferbeziehung
  - Alternativen
  - ...

## Optimierungsziel

$$\max \sum_{\text{Funktion } f} size_f \cdot n_f$$

In `lp_solve` -Syntax:

```
max : +64 n_f1 +48 n_f2 +42 n_f3 ;
```



# Flussbedingung: Initialer Aufruf

## Semantik

Der initiale Aufruf erfolgt maximal (wahlweise auch genau) ein mal

## Formalisierung

$$n_{\text{main}} \leq 1$$



## lp\_solve -Syntax

```
n_main <= 1;
```



# Flussbedingung: Mehrere Vorgänger

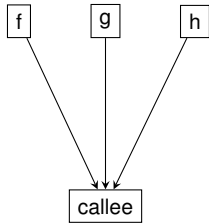
## Semantik

Jede Funktion kann nur so oft ausgeführt werden, wie sie von den Vorgängern aus aufgerufen wird

## Formalisierung

Sei  $f_{a \rightarrow b}$  die Anzahl der Aufrufe von b durch a:

$$n_{callee} \leq \sum_{p \in \text{Aufrufer}(callee)} f_{p \rightarrow callee}$$



## lp\_solve -Syntax

```
n_callee <= + f_f_callee + f_g_callee + f_h_callee ;
```



# Flussbedingung: Immer nur ein Nachfolger pro Funktion

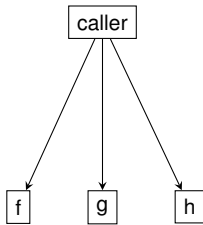
## Semantik

Jede Funktionsinkarnation ruft gleichzeitig jeweils maximal eine weitere Funktion auf

## Formalisierung

Sei  $f_{a \rightarrow b}$  die Anzahl der Aufrufe von  $b$  durch  $a$ :

$$\sum_{c \in \text{Aufgerufene}(\text{caller})} f_{\text{caller} \rightarrow c} \leq n_{\text{caller}}$$



## lp\_solve -Syntax

```
+ f_caller_f + f_caller_g + f_caller_h <= n_caller ;
```



# Flussbedingung: Rekursion

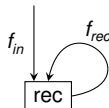
## Semantik

Rekursive Funktionen können pro Aufruf von außen bis zu ihrer maximalen Rekursionstiefe ( $d$ ) oft ausgeführt werden.

## Formalisierung

$$f_{rec} \leq d_{rec} \cdot f_{in}$$

$$n_{rec} \leq f_{in} + f_{rec}$$



## lp\_solve -Syntax

```
f_rec <= +42 f_in ;  
n_rec <= f_in + f_rec ;
```



- Problemformulierung in Ipsolve:

```
max: +40 n_main +20 n_f +60 n_g;
```

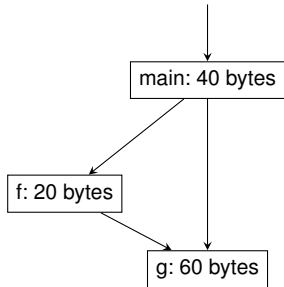
```
n_main <= 1;
```

```
+f_main_f +f_main_g <= n_main;
```

```
n_f <= +f_main_f;
```

```
+f_f_g <= n_f;
```

```
n_g <= +f_f_g +f_main_g;
```





# Beispiel

- Problemformulierung in lpsolve:

```
max: +40 n_main +20 n_f +60 n_g;
```

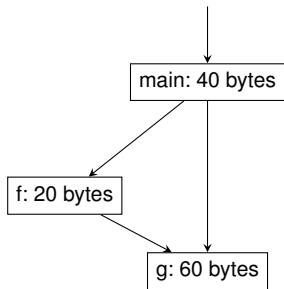
```
n_main <= 1;  
+f_main_f +f_main_g <= n_main;  
n_f <= +f_main_f;  
+f_f_g <= n_f;  
n_g <= +f_f_g +f_main_g;
```

- Ausgabe von lp\_solve :

```
Value of objective function: 120.00000000
```

```
Actual values of the variables:
```

```
n_main          1  
n_f             1  
n_g             1  
f_main_f       1  
f_main_g       0  
f_f_g          1
```



```
$ lp_solve infeasible.lp  
This problem is infeasible
```

### Infeasible Models

#### Logischer Widerspruch in Nebenbedingungen

Leider bietet `lp_solve` selbst direkt keine Hilfestellung zur Lokalisation. Die Entwickler empfehlen das Einführen von "slack"-Variablen:<sup>3</sup>

|                             |                                   |                     |
|-----------------------------|-----------------------------------|---------------------|
| <code>max: x + y;</code>    | <code>max: x + y</code>           | <code>x: 20</code>  |
| <code>x + 1 &lt;= x;</code> | <code>-1000 e_1</code>            | <code>y: 20</code>  |
| <code>y &gt; y + 1;</code>  | <code>-1000 e_2;</code>           | <code>e_1: 1</code> |
| <code>x &lt;= 20;</code>    | <code>x + 1 - e_1 &lt;= x;</code> | <code>e_2: 1</code> |
| <code>y &lt;= 20;</code>    | <code>y + e_2 &gt; y + 1;</code>  |                     |
|                             | <code>x &lt;= 20;</code>          |                     |
|                             | <code>y &lt;= 20;</code>          |                     |

<sup>3</sup><http://lpsolve.sourceforge.net/5.5/Infeasible.htm>

```
$ lp_solve unbounded.lp  
This problem is unbounded
```

## Unbounded Models

Eine oder mehrere der Variablen sind nach oben unbeschränkt

Durch künstliche Beschränkung aller Variablen im System (auf einen sehr großen Wert) lassen sich unbeschränkte Variablen detektieren:

|                    |                    |         |
|--------------------|--------------------|---------|
| max: $x + y + z$ ; | max: $x + y + z$ ; | x: 5000 |
| $z \leq y + 1$ ;   | $z \leq y + 1$ ;   | y: 20   |
| $y \leq 20$ ;      | $y \leq 20$ ;      | z: 21   |
|                    | $x \leq 5000$ ;    |         |
|                    | $y \leq 5000$ ;    |         |
|                    | $z \leq 5000$ ;    |         |



- `lp_solve` ist auf die Lösung linearer Gleichungssysteme ausgelegt
- Es ist dementsprechend nicht möglich, zwei Variablen zu multiplizieren
  - `a * b`  $\Rightarrow$  Syntaxfehler
  - `max : a b`  $\Rightarrow$  optimiert  $a + b$
- Lösung in VEZS für Konstanten (Stapelrahmengrößen): C-Präprozessor:

```
#define s_main 40  
#define s_f    20  
#define s_g    60
```

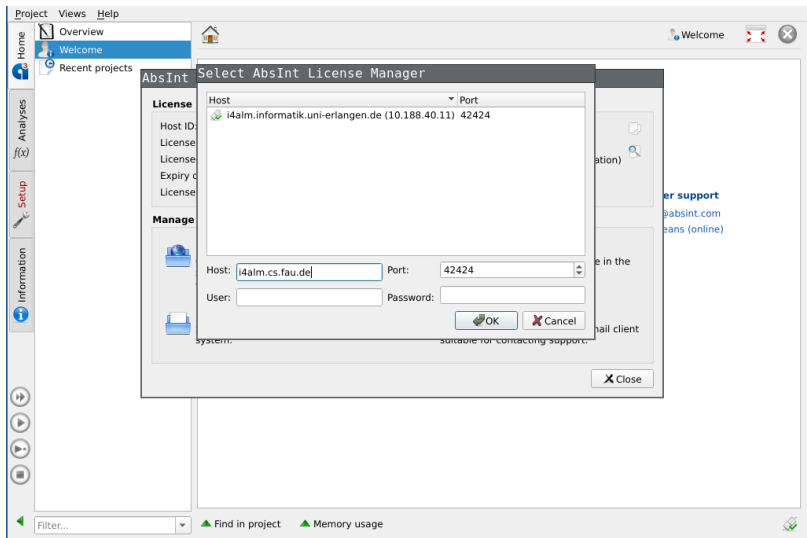
```
max: +s_main n_main +s_f n_f +s_g n_g;
```

~> `stackusage / lp_solvepp`



- Statische Code-Analyse mit a<sup>3</sup> Tool-Suite
  1. aiT: WCET-Analyse
  2. Stack-Analyzer: Stackbedarf
  3. ...
- Installiert im CIP-Pool
- `/proj/i4ezs/tools/a3_x86/bin/a3x86`

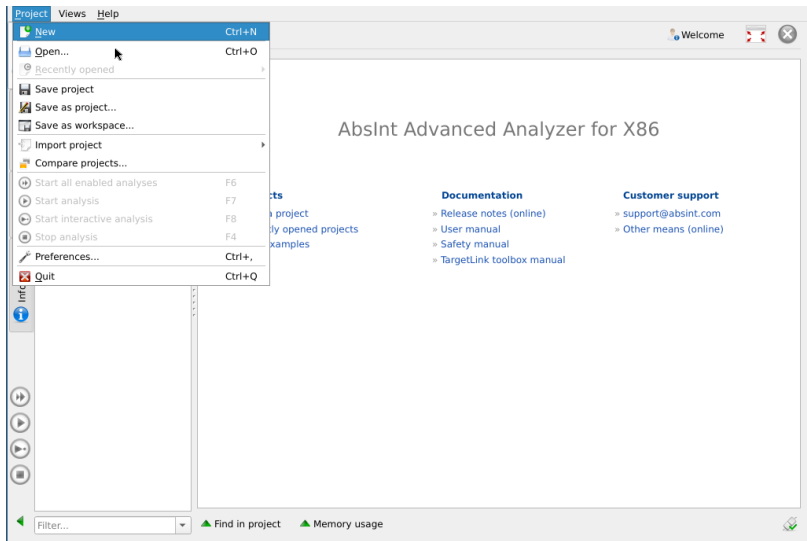




The screenshot shows the main window of the a<sup>3</sup> Analyzer software. The interface includes a menu bar (Project, Views, Help), a sidebar with navigation options (Home, Analyses, Setup, Information), and a main workspace. A dialog box titled "Select AbsInt License Manager" is open in the center. The dialog has a "License" section with a list of hosts, where "i4alm.informatik.uni-erlangen.de (10.188.40.11) 42424" is selected. Below this is a "Manage" section with input fields for "Host:" (containing "i4alm.cs.fau.de"), "Port:" (containing "42424"), "User:", and "Password:". A blue callout box with white text is overlaid on the dialog, containing the text "Zugangsdaten" and "Benutzer/Passwort wie bei RÜ-Helpdesk". To the right of the dialog, there is a sidebar with a search icon and a "Close" button. The background shows a "Welcome" screen with a "Recent projects" list and a "Filter..." dropdown at the bottom.

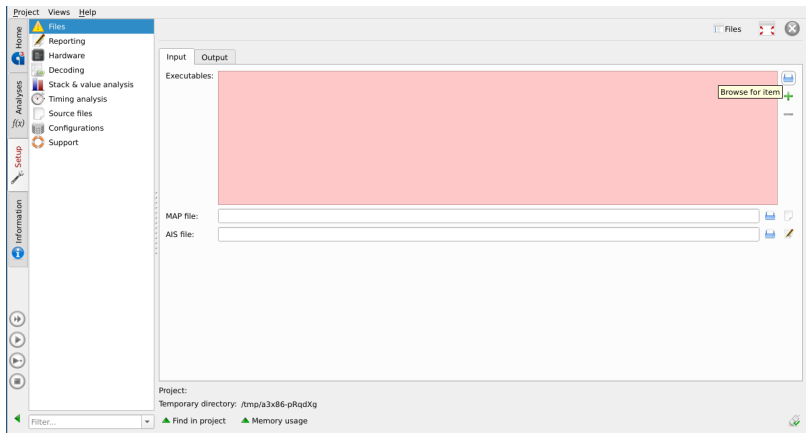


# a<sup>3</sup> Analyzer – Neues Projekt Anlegen



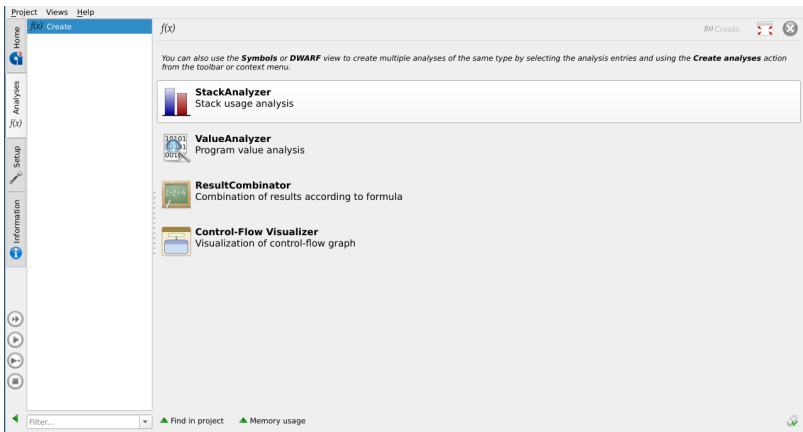


# a<sup>3</sup> Analyzer – Executable Angeben



The screenshot shows the 'Hardware' settings window in the a3 Analyzer. The window has a menu bar with 'Project', 'Hardware', 'Views', and 'Help'. On the left, there is a sidebar with categories: 'Home' (Files, Reporting, Hardware, Decoding), 'Analyses' (Stack & value analysis, Timing analysis, Source files, Configurations), 'Setup' (Support), and 'Information'. The main area is titled 'Machine Settings' and contains a 'General' tab. Under 'CPU', the 'Variant' is set to '64-bit (Long mode)'. Under 'Registers', there is a section for 'Register contents at start of analysis' with input fields for 'RSP', 'FS.base', 'GS.base', and 'Stack area'. The 'FS.base' and 'GS.base' fields have a 'hex' button next to them. At the bottom of the window, there is a search bar and buttons for 'Find in project' and 'Memory usage'.

# a<sup>3</sup> Analyzer – Stack-Analyse Selektieren



The screenshot shows the a3 Analyzer application window. The title bar reads "f(x) Create". The menu bar includes "Project", "Views", and "Help". On the left, a sidebar contains icons for "Home", "Analysis", "Setup", and "Information". The main area displays a list of analysis options:

- StackAnalyzer**: Stack usage analysis
- ValueAnalyzer**: Program value analysis
- ResultCombinator**: Combination of results according to formula
- Control-Flow Visualizer**: Visualization of control-flow graph

At the bottom of the window, there is a search bar labeled "Filter...", a "Find in project" button, and a "Memory usage" button. A small green checkmark icon is visible in the bottom right corner.



# a<sup>3</sup> Analyzer – Stack-Analyse Starten

The screenshot displays the a<sup>3</sup> Analyzer application window. The interface includes a menu bar (Project, Analysis, Views, Help) and a toolbar with icons for file operations and analysis. On the left, a sidebar contains navigation buttons: Home, Analysis, Setup, and Information. The main workspace is divided into two panes. The left pane shows a list of analysis items, with 'StackAnalyzer' selected. A yellow tooltip with the text 'Start all enabled analyses' is positioned over the 'Start' button in this pane. The right pane is titled 'StackAnalyzer' and contains configuration fields: ID (StackAnalyzer), Comment, Result (n/a), Configuration (Default Configuration), Dependencies (None), Analysis start (run), AIS file, and Expected result. There is also an unchecked checkbox for 'Enable ValueAnalyzer features'. At the bottom of the main workspace, there is a 'Filter...' dropdown and three status indicators: Messages, Find in project, and Memory usage.



# a<sup>3</sup> Analyzer – Analyseoutput

The screenshot displays the a3 Analyzer application window. The interface includes a sidebar with navigation options: Home, Analysis, f30, Setup, and Information. The main window title is "StackAnalyzer" and it shows the following details:

- ID:** StackAnalyzer
- Comment:** (empty)
- Result:** System: 96 bytes
- Settings:** Configuration: Default Configuration
- Dependencies:** None
- Analysis start:** run
- AIS file:** (empty)
- Expected result:** (empty)
- Enable ValueAnalyzer features

The main area displays the analysis output under the heading "Errors, warnings and info". The output is as follows:

```
StackAnalyzer - StackAnalyzer (0 Errors, 4 Warnings): Finished on 2020-06-15 at 17:10:14 after analyzing for 1 second
  Control-Flow & Stack Analysis
    Reading binary 'stacktest'.
    #1039: ELF file is not an executable, but shared object file.
    #1033: ELF file is not a statically linked executable, but contains relocations.
    #1034: ELF file is not a statically linked executable, but contains dynamic link information.
    Using decoder for 'x86_64' and compiler 'GCC'.
    Recursion 0x1125 'h' found, recursion members:
    Value analyzer statistics (max length=2, default-unroll=2, normal mode):
    Loop analysis found 0 loop bounds.
    #1009: For routine 'h' the default incarnation limit of 1 is used.
    The analyzer optimized the stack graph of entry 'run' from 5/5 to 2/2 nodes * calls (non-optimizable routines: 2).
    The analyzer optimized the stack graph of entry 'run' from 5/5 to 4/4 nodes * calls (non-optimizable routines: 2).
    Maximum global stack height: 96
    Last process took 0 s and used not more than 30 MB (RSS 13 MB) of memory
  Reporting
  Creating HTML report
  Finished on 2020-06-15 at 17:10:14 after analyzing for 1 second with 0 errors, 4 warnings
```

At the bottom of the window, there is a filter input field, buttons for "Messages", "Find in project", and "Memory usage", and the text "Overall analysis time: <1s".

# a<sup>3</sup> Analyzer – Analyseoutput

The screenshot shows the StackAnalyzer application window. The main panel displays the analysis results for a project named 'StackAnalyzer'. The results are categorized into 'Errors, warnings and info' and 'Latest log'. A yellow warning message is highlighted, indicating a warning to ignore ELF files. The warning text is as follows:

- Control-Flow & Stack Analysis**
- #1039: ELF file is not an executable, but shared object file.**
- #1033: ELF file is not a statically linked executable, but contains relocations.**
- #1034: ELF file is not a statically linked executable, but contains dynamic link information.**
- Using decoder for 'x86\_64' and compiler 'GCC'.
- Recursion 0x1125 'h' found, recursion members:
- Value analyzer statistics (max length=2, default-unroll=2, normal mode):
- Loop analysis found 0 loop bounds.
- The analyzer optimized the stack graph of entry 'run' from 5/5 to 2/2 nodes.
- #1097: For routine 'h' the default incarnation limit of 1 is used.**
- The analyzer optimized the stack graph of entry 'run' from 5/5 to 4/4 nodes \* calls (non-optimizable routines: 2).
- Maximum global stack height: 96
- Last process took 0 s and used not more than 30 MB (RSS 13 MB) of memory
- Reporting
- Creating HTML report
- Finished on 2020-06-15 at 17:10:14 after analyzing for 1 second with 0 errors, 4 warnings

A pink callout box with an arrow points to the warning messages, containing the text: **⇒ Warnung zu ELF ignorieren**



# a<sup>3</sup> Analyzer – Callgraph

The screenshot displays the a3 Analyzer interface for the StackAnalyzer tool. The top menu bar includes 'Project', 'Analysis', 'Views', and 'Help'. The left sidebar shows navigation options: 'Home', 'Analysis', 'f30', 'Setup', and 'Information'. The main window title is 'StackAnalyzer' and it shows the tool's ID, comment, and result (System: 96 bytes). The 'Settings' tab is active, showing configuration options like 'Default Configuration', 'Dependencies: None', and 'Analysis start: run'. Below the settings is a list of messages. The message '#21027 Recursion in the control flow' is selected, and a context menu is open over it, listing actions such as 'Copy', 'Copy part', 'Show in call graph', 'Show in disassembly', 'Show in file', 'Copy path', 'Copy AIS annotation', 'Find limit in DWARF', 'Show all folded messages of this type', 'Show all folded messages', 'Reset state of all folded messages', 'Clear all', 'Expand recursively', 'Collapse recursively', 'Expand all', and 'Collapse all'. The bottom status bar indicates 'Overall analysis time: <1s'.



# a<sup>3</sup> Analyzer – Annotationstemplate kopieren

The screenshot displays the a3 Analyzer interface. At the top, a title bar reads "Project Analysis Graph Views Help". Below it, a toolbar contains icons for file operations and analysis settings. The main window is titled "Analysis graph" and shows a stack graph with three nodes: "run: [-B..32]", "g: [-B..32]", and "[REC...". A context menu is open over the "[REC..." node, listing various actions such as "Toggle fold", "Show address in disassembly", "Copy AIS annotation", and "Recursion bounds".

Below the graph, a panel titled "Errors, warnings and info" displays a list of warnings. The first warning is highlighted in yellow:

- StackAnalyzer - StackAnalyzer (0 Errors, 4 Warnings): Finished on 2020-06-15 at 17:10:14**
  - Control-Flow & Stack Analysis
    - Reading binary 'stacktest'.
      - #1039: ELF file is not an executable, but shared object file.
      - #1033: ELF file is not a statically linked executable, but contains relocations.
      - #1034: ELF file is not a statically linked executable, but contains dynamic link information using decoder for 'x86\_64' and compiler 'GCC'.
        - Recursion 0x1125 'h' found, recursion members:
          - Value analyzer statistics (max-length=2, default-unroll=2, normal mode):
            - Loop analysis found 0 loop bounds.
  - Stack graph
    - #1077: For routine 'h' the default incrementation limit of 1 is used.
      - The analyzer optimized the stack graph of entry 'run' from 5/5 to 2/2 nodes \* calls (non-optimized)
      - The analyzer optimized the stack graph of entry 'run' from 3/5 to 4/4 nodes \* calls (non-optimized)
      - Maximum global stack height: 96
  - Reporting
    - Creating HTML report
    - Finished on 2020-06-15 at 17:10:14 after analyzing for 1 second with 0 errors, 4 warnings





# a<sup>3</sup> Analyzer – Stack-Analyse Starten

The screenshot shows the StackAnalyzer application window. The main area is titled 'StackAnalyzer' and contains the following fields and controls:

- ID:** StackAnalyzer
- Comment:** (empty text box)
- Result:** n/a
- Settings** (selected) / **Output** (tab)
- Configuration:** Default Configuration (dropdown menu)
- Dependencies:** None (with a refresh icon)
- Analysis start:** run (with a refresh icon)
- AIS file:** (empty text box with a file selection icon)
- Expected result:** (empty text box)
- Enable ValueAnalyzer feature**

An orange callout box points to the 'Enable ValueAnalyzer feature' checkbox with the text: **⇒ .ais-Datei für benutzerdefinierte Annotationen**

On the left sidebar, the 'Start all enabled analyses' button is highlighted with a yellow box.



# a<sup>3</sup> Analyzer – Annotationstemplate kopieren

The screenshot displays the a3 Analyzer's Analysis graph window. At the top, a green box indicates "Maximum Stack Usage for Entry 'run': 96". Below this, a control-flow graph shows two nodes: "run: [-B..32]" and "g: [-B..32]". A context menu is open over the "g" node, listing various actions such as "Toggle fold", "Show address in disassembly", "Copy AIS annotation", and "Recursion bounds". The "Recursion bounds" sub-menu is also visible, showing options like "Incarnation limit", "Enter with", "Infeasible", and "Not analyzed".

Errors, warnings and info | Latest log

**StackAnalyzer – StackAnalyzer (0 Errors, 4 Warnings): Finished on 2020-06-15 at 17:10:14**

- Control-Flow & Stack Analysis
  - Reading binary 'stacktest'.
    - #1039: ELF file is not an executable, but shared object file.
    - #1033: ELF file is not a statically linked executable, but contains relocations.
    - #1034: ELF file is not a statically linked executable, but contains dynamic link information using decoder for 'x86\_64' and compiler 'GCC'.
      - Recursion 0x1125 'h' found, recursion members:
        - Value analyzer statistics (max length=2, default-unroll=2, normal mode):
          - Loop analysis found 0 loop bounds.
      - The analyzer optimized the stack graph of entry 'run' from 5/5 to 2/2 nodes \* calls (non-optimized)
    - #1077: For routine "h" the default incarnation limit of 1 is used.
      - The analyzer optimized the stack graph of entry 'run' from 5/5 to 4/4 nodes \* calls (non-optimized)
      - Maximum global stack height: 96
      - Last process took 0 s and used not more than 30 MB (RSS 13 MB) of memory
- Reporting
  - Creating HTML report
  - Finished on 2020-06-15 at 17:10:14 after analyzing for 1 second with 0 errors, 4 warnings

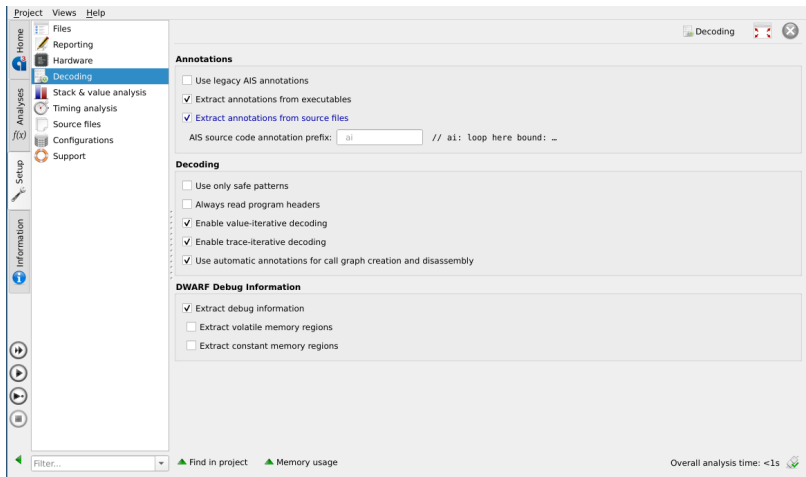
Overall analysis time: <1s

## Ais-Notationen

- Auch als C-Kommentar verwendbar
- // ai: routine "h" recursion bound : 0 .. 42;



# a<sup>3</sup> Analyzer – Kommentar-Parsing Aktivieren



The screenshot shows the 'Decoding' settings window in the a<sup>3</sup> Analyzer. The window has a title bar with 'Decoding' and standard window controls. On the left is a sidebar with categories: Home, Analyses (with 'Decoding' selected), Setup, and Information. The main area is divided into three sections: Annotations, Decoding, and DWARF Debug Information. At the bottom, there are status indicators for 'Find in project' and 'Memory usage', and a timer showing 'Overall analysis time: <1s'.

**Annotations**

- Use legacy AIS annotations
- Extract annotations from executables
- Extract annotations from source files

AIS source code annotation prefix:  // ai: loop here bound: \_

**Decoding**

- Use only safe patterns
- Always read program headers
- Enable value-iterative decoding
- Enable trace-iterative decoding
- Use automatic annotations for call graph creation and disassembly

**DWARF Debug Information**

- Extract debug information
- Extract volatile memory regions
- Extract constant memory regions

Filter... Find in project Memory usage Overall analysis time: <1s



- Existierende Implementierung: Array-Datenstruktur
- Vorgegebene Funktionen: Sortieren, Maximumssuche, ...
- Aufgaben
  1. Dynamische Analyse
    - 1.1 Thread erstellen
    - 1.2 Stack initialisieren
    - 1.3 Programm (mit Eingabedaten) ausführen
    - 1.4 Stackverbrauch messen
  2. Statische Analyse
    - 2.1 ILP aus Aufrufgraph aufstellen
    - 2.2 Mittels `lp_solve` lösen
    - 2.3 Analyse mittels `a3` Stack-Analyzer

