# Seminar Ausgewählte Kapitel der Systemsoftware (AKSS)

Bluetooth Low Energy In Intermittently-Powered Wireless Communication Systems - An Introduction

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- 1. Introduction
- 2. Bluetooth Low Energy (BLE)
- 3. Intermittent Implementation FreeBie
- 4. Findings & Conclusion

# Introduction

## Introduction

- Bluetooth Low Energy (BLE) used in many IoT sensors & devices
- BLE Real-time requirements
- Enable undisrupted bi-directional communication in intermittent systems
- Battery-free, using capacitors and energy harvesting



# Bluetooth Low Energy (BLE)

- Introduced in Bluetooth version 4.0 (2010)
- Low power variant of Bluetooth Classic
- Different protocol stack and hardware
- Bursty and low bandwidth communication
- Able to work on coin-cell batteries
- $\Rightarrow$  Incompatible because fundamentally different
  - But most modern consumer devices support both

- BLE operates in the ISM Band (2.4 GHz electromagnetic frequency spectrum)
- Shared with WLAN, medical instruments, microwaves, ...
- Spectrum subdivided into 40 channels
  - 3 advertising channels for pairing
  - 37 data channels for transmissions
  - Adaptive frequency hopping algorithm in data channels to avoid noise & other signals
- Frequency ("key") for "0" and "1" in each channel
- ⇒ Gaussian Frequency Shift Keying for data transfer

### **BLE Connection Phases**



https://www.bluetooth.com/blog/bluetooth-pairing-part-1-pairing-feature-exchange/

previous attempts:

- too many undesirable traits
- poor performance, lengthy reconnection (around 40s)
- breaking protocol's timing and real-time requirements



# Intermittent Implementation -FreeBie

# **FreeBie Architecture**

- A BLE ARM-based microcontroller unit (MCU)
- B Non-volatile FRAM
- C external Real-Time Clock (RTC)
- **D** Capacitors
- E Energy harvester
- solar panels & display on the back





J. de Winkel, H. Tang, and P. Pawełczak. Intermittently-Powered Bluetooth That Works. In Proceedings of the 20th Annual International Conference on Mobile Systems, Applications and Services. MobiSys '22, pp. 287-301.

- 1. Time-Deterministic Checkpoint and Real-Time Restoration
- 2. Virtualisation Layer
- 3. Dynamic Network Handling

# Time-Deterministic Checkpoint and Real-time Restoration

- Checkpoint is a function that is introduced to the original protocol
- Saves system state until the function call
- Triggered by end of wireless protocol process
- Stored in non-volatile FRAM
- Process & memory separation into network, application & OS
- Classify processes as real-time or non-real-time
- Scheduler decides whether to sleep or turn off power to MCU

- 1. Determine next power-on time  $T_{wakeUp}$
- 2. Determine real-time processes for restoration
- 3. Checkpoint processes
- 4. OS Checkpoint
- 5. Switch off power domain

#### Real-time processes:

- Restored at boot, before scheduler starts
- If checkpoint present:
  - restore OS and network
  - restore real-time application processes
- Synchronise to external RTC
- Compensate for power-off time
- Resume normal operation

#### Non-real-time processes:

Restored dynamically prior to execution

## Checkpoint, Restoration, Memory Map Overview

#### FRAM non-volatile, slower



SRAM volatile, faster

Masking time and intermittent effects to host and applications to support time-deterministic restoration.

- System is divided into processor power domain & "always-on" ultra-low-power domain
- Energy consuming components like processor can be switched off, if next process is scheduled later than T<sub>minOff</sub> = 20ms
- External RTC belongs to always-on domain
- $\Rightarrow$  RTC can be set to switch processor on again at  $T_{wakeUp}$

#### T<sub>wakeUp</sub>

 $T_{wakeUp} = T_{sync} - T_{startUp} - T_{restore}$ 

## $T_{startUp}$ = 10 ms (start up duration of system)

- $T_{restore}$  = 10 ms
- T<sub>sync</sub> := known time of next sync with external RTC needed to reset internal MCU time registers after power failure to correct time (from external RTC) ⇒ masking intermittent time effects

#### Main areas:

- 1 Network recovery
- 2 Dynamic network adaptation

Improves:

- System connectivity
- Performance
- Energy efficiency

#### $\textbf{Parameters} \rightarrow \text{can} \text{ be adjusted}$

```
Supervision Timeout (ST) [s] \in [0, 32]
```

Time after which a connection is considered lost (i.e. Connection Timeout)

## Connection Interval (CI) [s] $\in$ [1, 4]

Time between two consecutive data transfers between devices

```
Slave Latency (SL) [int] \in [0, 3]
```

How many connection events can be skipped by end device

Needed when system ...

- Performs normal restoration (after power failure or power off)
- Was unable to turn back on at T<sub>wakeUp</sub>
- $\Rightarrow$  Possible when still within ST window
- $\Rightarrow$  ST set to maximum of 32s

Steps:

- 1. Skip missed connection events
- 2. Schedule network process for next connection event
- 3. Retransmit lost packages (or attempt to)

Adjusting CI and SL depending on energy levels:

Energy Level (Luminosity)	CI	SL	ST
Low (200-300lx)	4S	3	32S
Medium (600lx)	25	1	32S
High (10klx)	1S	0	32S

- Low energy: increase CI and SL
  - Lower throughput and bandwidth
  - $\Rightarrow$  But energy available longer
- High energy: decrease CI and SL
  - ⇒ Increases responsiveness, throughput and overall performance

# **Findings & Conclusion**

## Findings & Conclusion (1-2 slides)

- FreeBie evaluation mostly successful
- Most notably 24h benchmark & connection retention
- In idle up to 9.5 times less energy consumption than reference device



FreeBie network activity (blue) around power failure (pink) and connection retention

# **Findings & Conclusion**

#### But significant external memory access overhead



# FreeBie-C: memory overhead data removed, showing significant impact

⇒ different MCU with internal non-volatile FRAM/MRAM could be used to eliminate external memory

# **Questions?**