

vNV-Heap: An Ownership-Based Virtually Non-Volatile Heap for Embedded Systems

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Friedrich-Alexander-Universität
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Saarland Informatics
Campus



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Energy-Harvesting Systems with Intermittent Power Supply

Observation Devices

Remote environmental sensor [1]

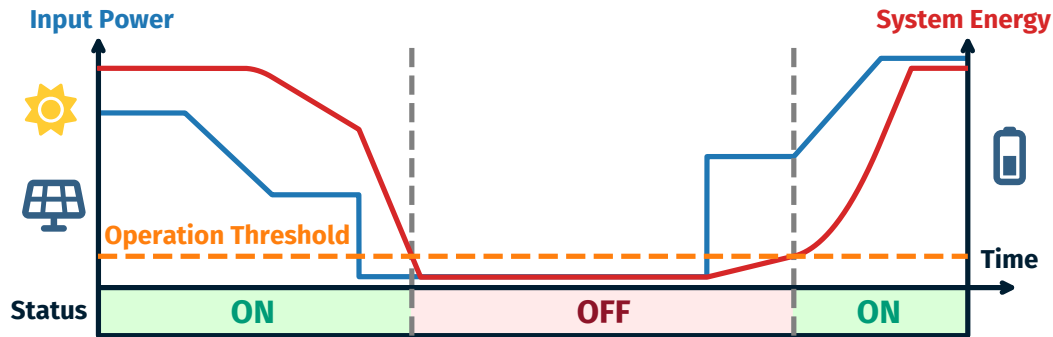


IoT Devices

Battery-free BLE smartwatch [3]



Intermittent Computing



- **Hard intermittency:** Device shuts down when $energy < threshold$
- **State retention:** Persist state in time! ($\rightarrow$ predictability)

Large Amounts of State

- **Memory-intense applications**
(→ DNNs, LLMs)
 - **Store work packages to be processed later**
(→ sensor data, IP packets)
- ⇒ **Commonly:** Explicitly allocated memory (→ heaps)



State Retention Mechanism

Q *How do you persist state before the system shuts down?*

A **Copy data from volatile RAM to NVM**

Q *What data should you persist?*

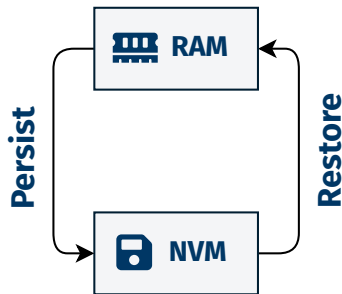
A ~~Persist all data~~

A **Persist only modified data**

A **Limit amount of modified state**

Takeaway #1

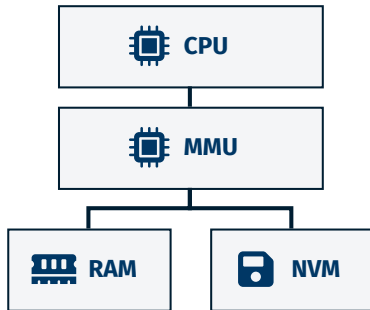
Precise modification tracking is required!



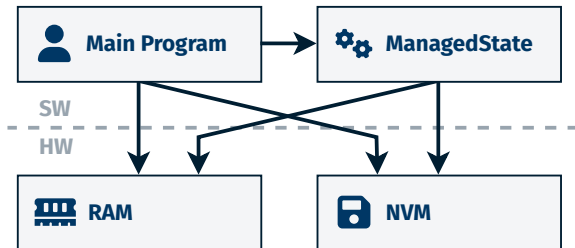
MMU-Based Approaches

1. NVM safety?
2. Transparent swap?
3. Predictable?
4. HW agnostic?

MMU-based



ManagedState [2] as SW-Only Approach



Modifying Data

```
1  modify_start(ptr);
2  *ptr += 1; // valid
3  modify_end(ptr);
4
5  // integrity at risk
6  *ptr += 1;
```

[2] S. Sliper et al. “Efficient State Retention through Paged Memory Management for Reactive Transient Computing”. In: DAC '19. ACM, June 2019

Comparison of Existing Approaches

	MMU-based	ManagedState	<u>vNV-Heap</u>
1. NVM safety?	✓	✗	✓
2. Transparent swap?	✓	✗	✓
3. Predictable?	✗	✓	✓
4. HW agnostic?	✗	✓	✓

vNV-Heap: virtually Non-Volatile Hheap

Comparison of Existing Approaches

	MMU-based	ManagedState	<u>vNV-Heap</u>
1. NVM safety?	✓	✗	✓
2. Transparent swap?	✓	✗	✓
3. How does vNV-Heap achieve this? Pre-allocated	✗	✓	✓
4. HW agnostic?	✗	✓	✓

vNV-Heap: virtually Non-Volatile Heap

Comparison of Existing Approaches

	MMU-based	ManagedState	<u>vNV-Heap</u>
1. NVM safety?	✓	✗	✓
2. Transparent swap?	✓	✗	✓
3. Predictable?	✗	✗	✓
4. HW agnostic?	✗	✓	✓

Idea: Leverage Modern Programming Languages!

vNV-Heap: virtually Non-Volatile Heap



Zero-cost abstraction: Statically checked



Ownership

- Each value has an owner
- There exists only one owner
- The value is dropped once the owner goes out of scope

Borrowing

There exists either:

- $= 1$ mutable reference or
- ≥ 0 immutable references

Memory Management using Ownership & Borrowing

Previous Example: Access and modify value

(a) Unsafe

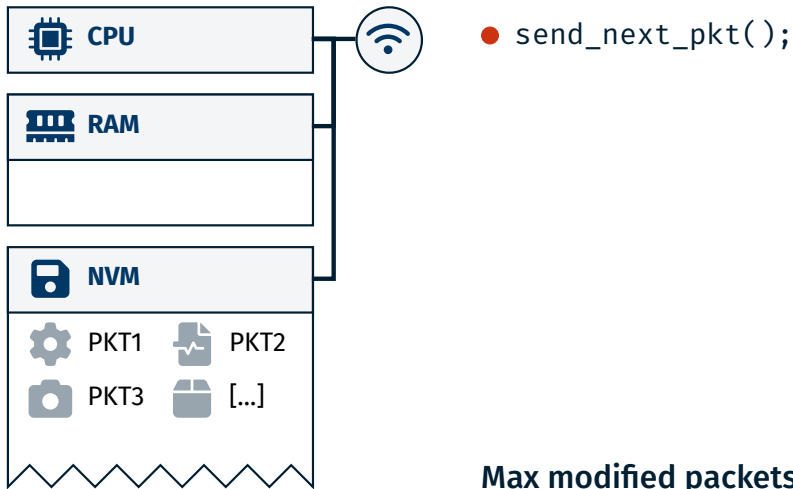
```
1  modify_start(ptr);
2  *ptr += 1; // valid
3  modify_end(ptr);
4
5  // integrity at risk
6  *ptr += 1;
```

(b) Ownership-Checked

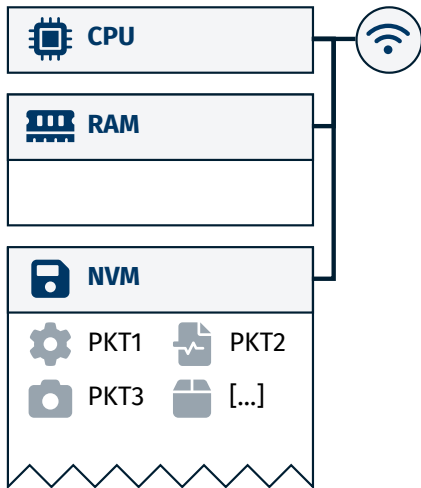
```
1  ref = get_mut(handle);
2  *ref += 1; // valid
3  drop(ref);
4
5  // compilation error
6  *ref += 1;
```

⇒ **Overall:** Makes SW-based modification tracking & swapping safe!

Practical Example



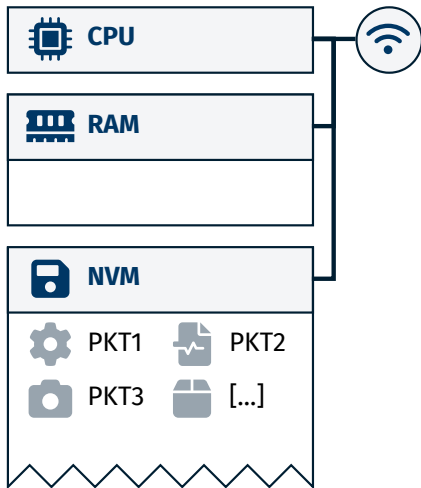
Practical Example



```
send_next_pkt();  
● → pkt1 = dequeue_pkt();
```

Max modified packets: 1

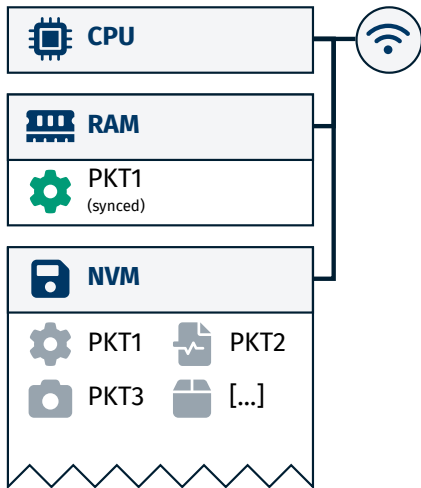
Practical Example



```
send_next_pkt();  
→ pkt1 = dequeue_pkt();  
● → ref1 = pkt1.get_mut();
```

Max modified packets: 1

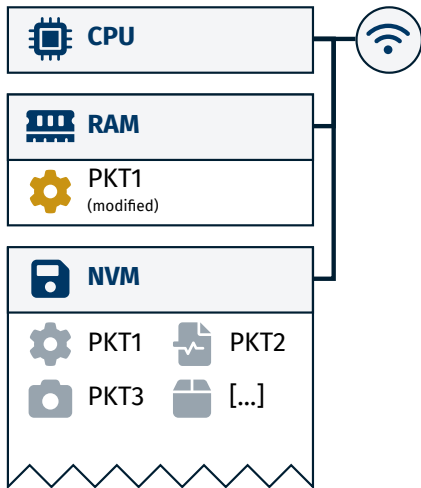
Practical Example



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send_next_pkt();  
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Max modified packets: 1

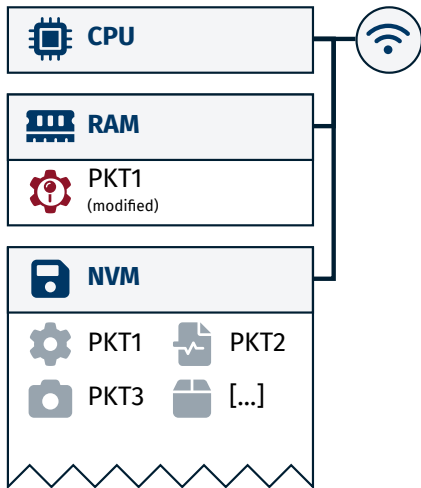
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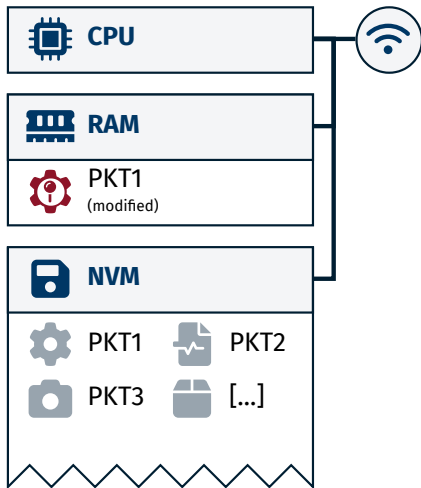
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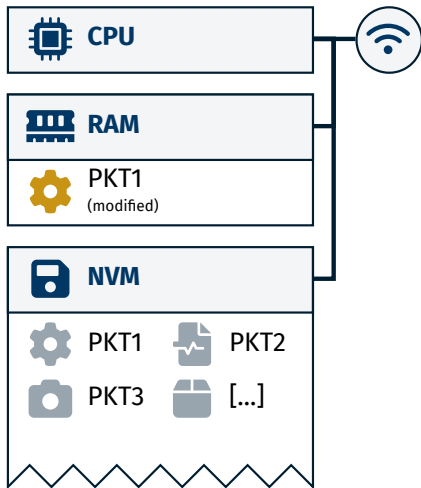
Practical Example



```
send_next_pkt();  
→ pkt1 = dequeue_pkt();  
→ ref1 = pkt1.get_mut();  
● → ref1.crc = calc_crc(ref1);
```

Max modified packets: 1

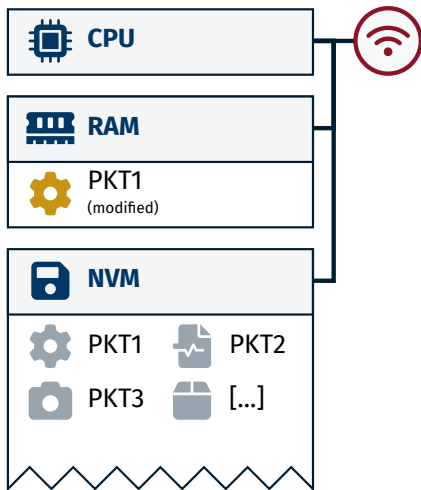
Practical Example



```
send_next_pkt();  
→ pkt1 = dequeue_pkt();  
→ ref1 = pkt1.get_mut();  
→ ref1.crc = calc_crc(ref1);  
● → drop(ref1);
```

Max modified packets: 1

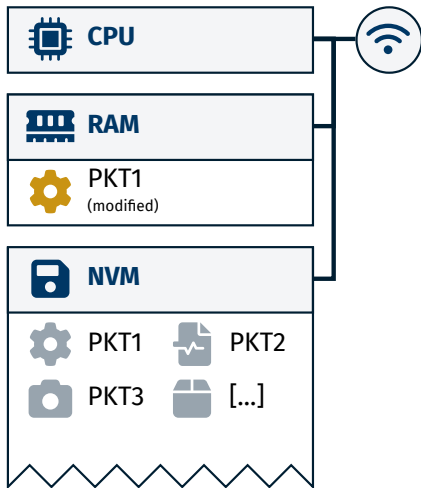
Practical Example



```
send_next_pkt();  
→ pkt1 = dequeue_pkt();  
→ ref1 = pkt1.get_mut();  
→ ref1.crc = calc_crc(ref1);  
→ drop(ref1);  
● → send(pkt1);
```

Max modified packets: 1

Practical Example



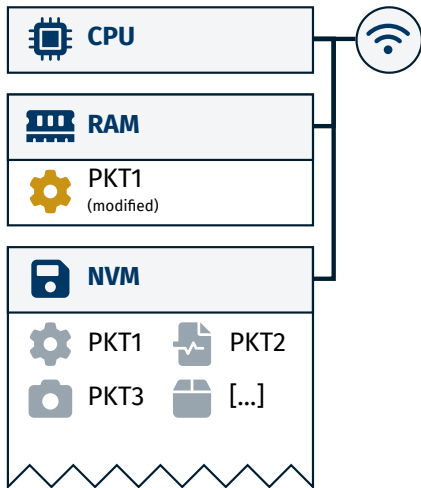
```
send_next_pkt();
```

```
→ [...]
```

● `send_next_pkt();`

Max modified packets: 1

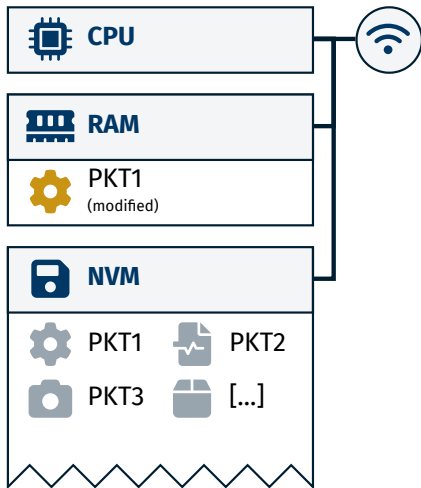
Practical Example



```
send_next_pkt();  
→ [...]  
send_next_pkt();  
● → pkt2 = dequeue_pkt();
```

Max modified packets: 1

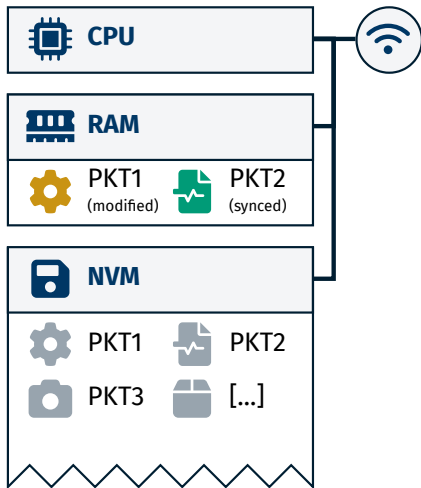
Practical Example



```
send_next_pkt();  
→ [...]  
send_next_pkt();  
→ pkt2 = dequeue_pkt();  
● → ref2 = pkt2.get_mut();
```

Max modified packets: 1

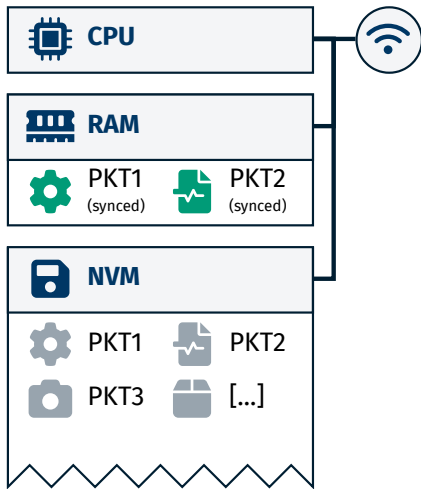
Practical Example



```
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→ [...]  
send_next_pkt();  
→ pkt2 = dequeue_pkt();  
● → ref2 = pkt2.get_mut();
```

Max modified packets: 1

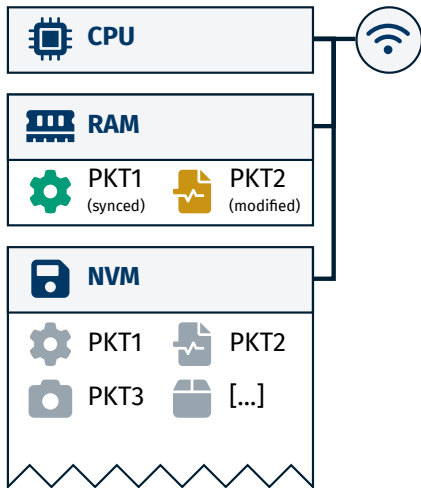
Practical Example



```
send_next_pkt();  
→ [...]  
send_next_pkt();  
→ pkt2 = dequeue_pkt();  
● → ref2 = pkt2.get_mut();
```

Max modified packets: 1

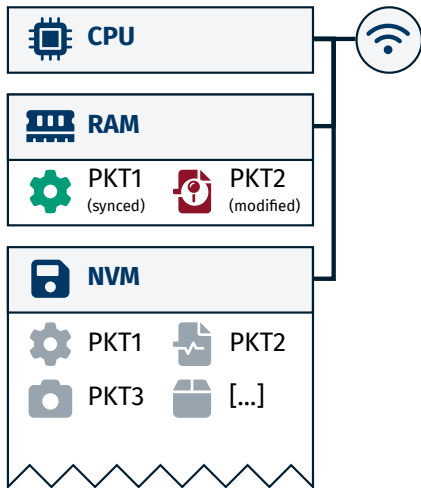
Practical Example



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→ [...]  
send_next_pkt();  
→ pkt2 = dequeue_pkt();  
● → ref2 = pkt2.get_mut();
```

Max modified packets: 1

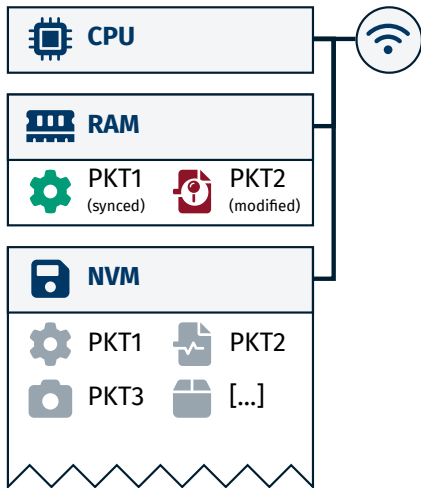
Practical Example



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send_next_pkt();  
→ pkt2 = dequeue_pkt();  
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```

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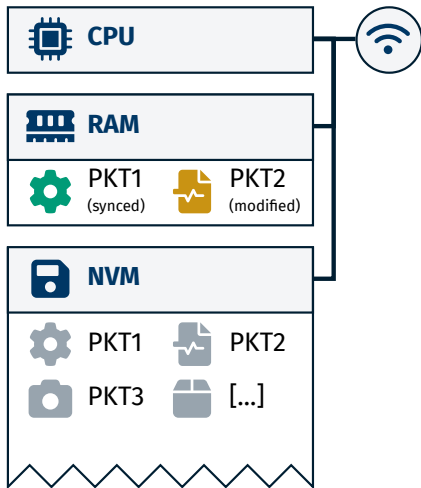
Practical Example



```
send_next_pkt();  
→ [...]  
send_next_pkt();  
→ pkt2 = dequeue_pkt();  
→ ref2 = pkt2.get_mut();  
● → ref2.crc = calc_crc(ref2);
```

Max modified packets: 1

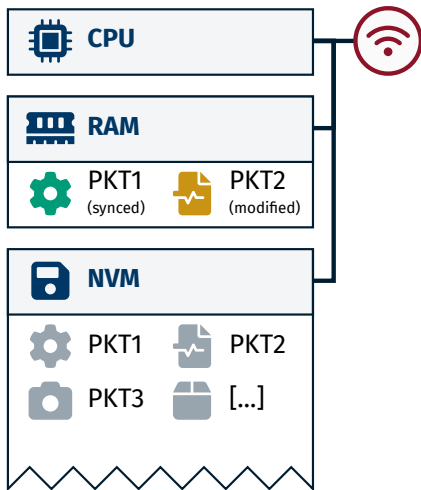
Practical Example



```
send_next_pkt();  
→ [...]  
send_next_pkt();  
→ pkt2 = dequeue_pkt();  
→ ref2 = pkt2.get_mut();  
→ ref2.crc = calc_crc(ref2);  
● → drop(ref2);
```

Max modified packets: 1

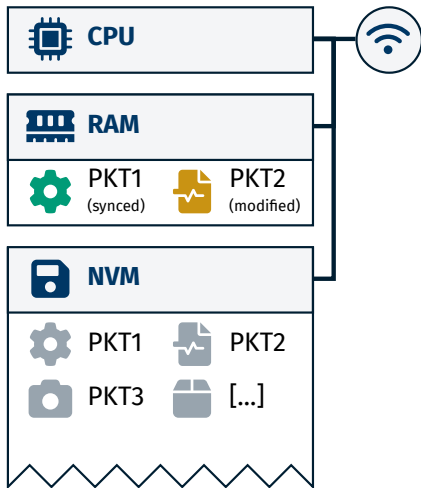
Practical Example



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→ [...]  
send_next_pkt();  
→ pkt2 = dequeue_pkt();  
→ ref2 = pkt2.get_mut();  
→ ref2.crc = calc_crc(ref2);  
→ drop(ref2);  
● → send(pkt2);
```

Max modified packets: 1

Practical Example



```
send_next_pkt();
```

```
→ [...]
```

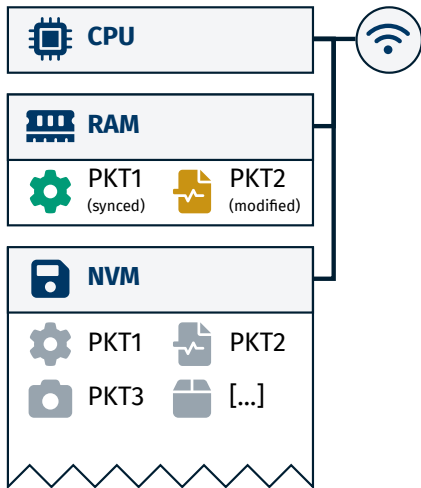
```
send_next_pkt();
```

```
→ [...]
```

```
● send_next_pkt();
```

Max modified packets: 1

Practical Example



```
send_next_pkt();
```

```
→ [...]
```

```
send_next_pkt();
```

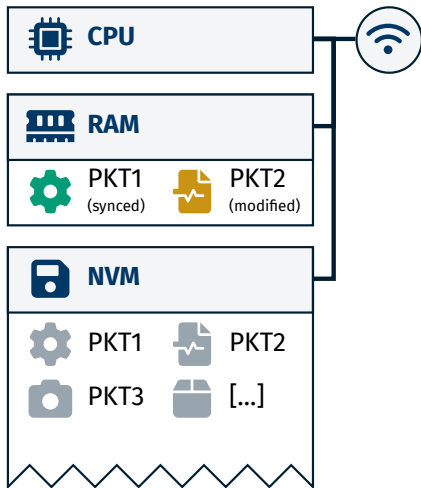
```
→ [...]
```

```
send_next_pkt();
```

```
● → pkt3 = dequeue_pkt();
```

Max modified packets: 1

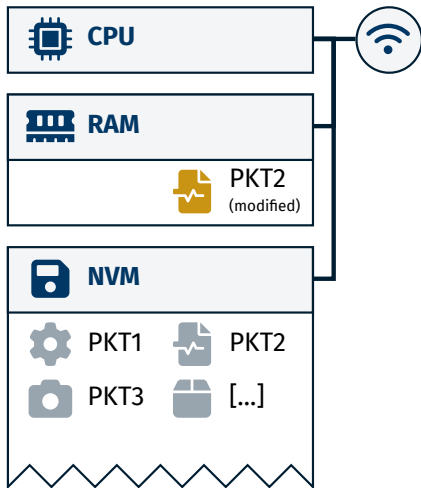
Practical Example



```
send_next_pkt();  
→ [...]  
send_next_pkt();  
→ [...]  
send_next_pkt();  
→ pkt3 = dequeue_pkt();  
● → ref3 = pkt3.get_mut();
```

Max modified packets: 1

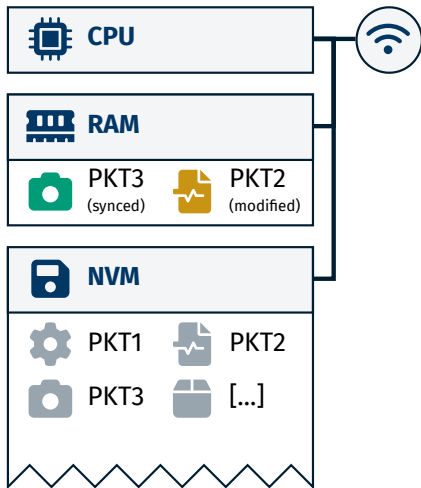
Practical Example



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→ [...]  
send_next_pkt();  
→ [...]  
send_next_pkt();  
→ pkt3 = dequeue_pkt();  
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```

Max modified packets: 1

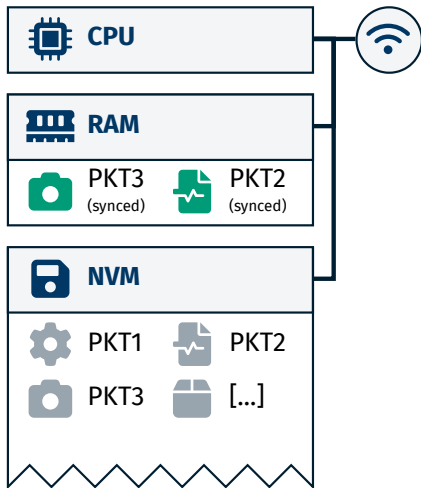
Practical Example



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→ [...]  
send_next_pkt();  
→ [...]  
send_next_pkt();  
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● → ref3 = pkt3.get_mut();
```

Max modified packets: 1

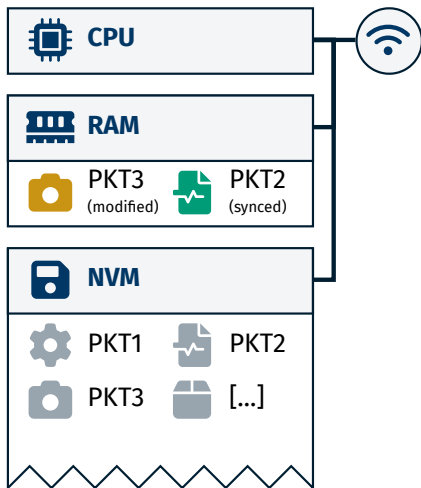
Practical Example



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→ [...]  
send_next_pkt();  
→ [...]  
send_next_pkt();  
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```

Max modified packets: 1

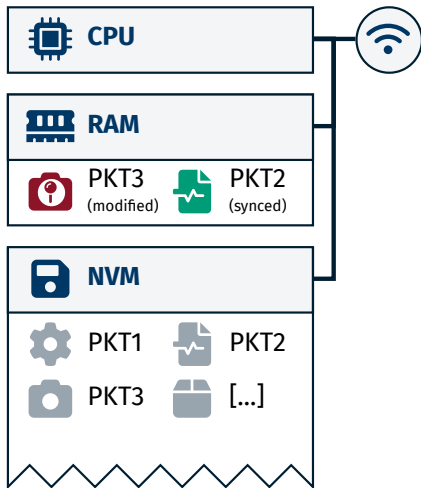
Practical Example



```
send_next_pkt();  
→ [...]  
send_next_pkt();  
→ [...]  
send_next_pkt();  
→ pkt3 = dequeue_pkt();  
● → ref3 = pkt3.get_mut();
```

Max modified packets: 1

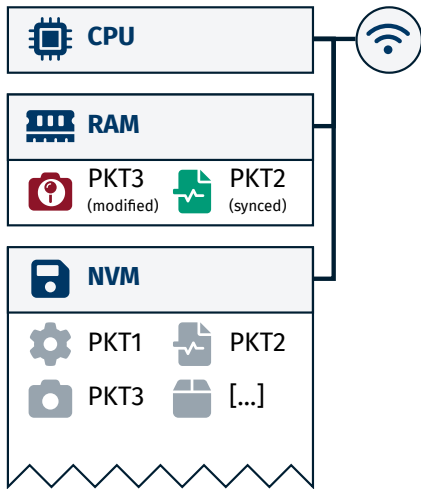
Practical Example



```
send_next_pkt();  
→ [...]  
send_next_pkt();  
→ [...]  
send_next_pkt();  
→ pkt3 = dequeue_pkt();  
● → ref3 = pkt3.get_mut();
```

Max modified packets: 1

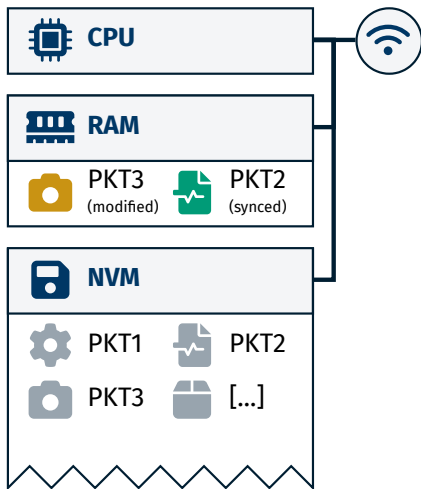
Practical Example



```
send_next_pkt();  
→ [...]  
send_next_pkt();  
→ [...]  
send_next_pkt();  
→ pkt3 = dequeue_pkt();  
→ ref3 = pkt3.get_mut();  
● → ref3.crc = calc_crc(ref3);
```

Max modified packets: 1

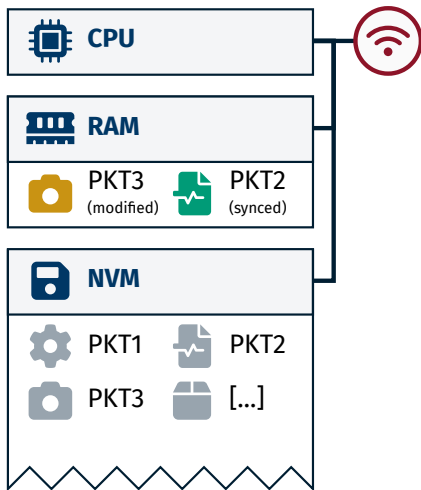
Practical Example



```
send_next_pkt();  
→ [...]  
send_next_pkt();  
→ [...]  
send_next_pkt();  
→ pkt3 = dequeue_pkt();  
→ ref3 = pkt3.get_mut();  
→ ref3.crc = calc_crc(ref3);  
● → drop(ref3);
```

Max modified packets: 1

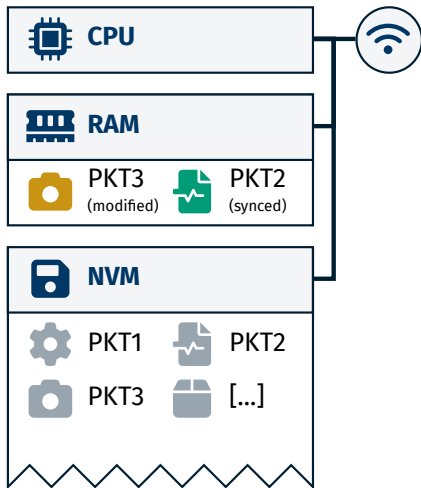
Practical Example



```
send_next_pkt();  
→ [...]  
send_next_pkt();  
→ [...]  
send_next_pkt();  
→ pkt3 = dequeue_pkt();  
→ ref3 = pkt3.get_mut();  
→ ref3.crc = calc_crc(ref3);  
→ drop(ref3);  
● → send(pkt3);
```

Max modified packets: 1

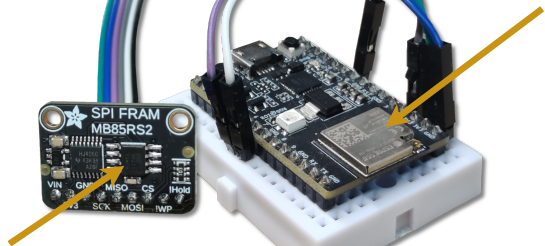
Practical Example



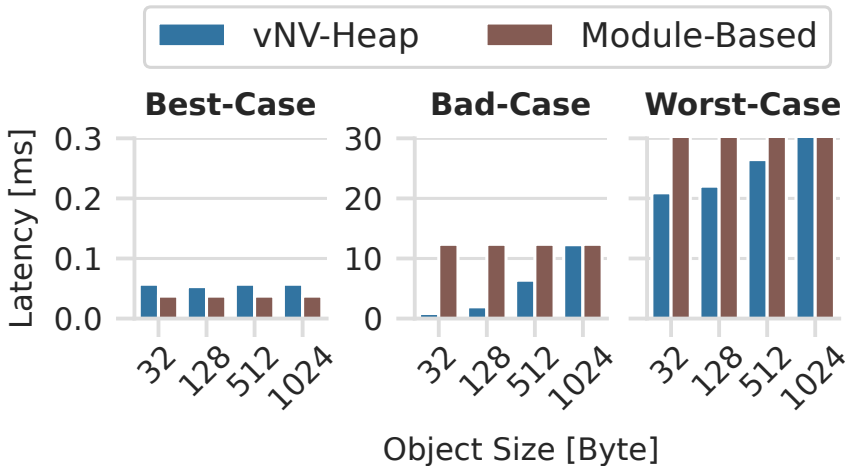
```
send_next_pkt();  
→ [...]  
send_next_pkt();  
→ [...]  
send_next_pkt();  
→ pkt3 = dequeue_pkt();  
→ ref3 = pkt3.get_mut();  
→ ref3.crc = calc_crc(ref3);  
→ drop(ref3);  
→ send(pkt3);
```

Max modified packets: 1

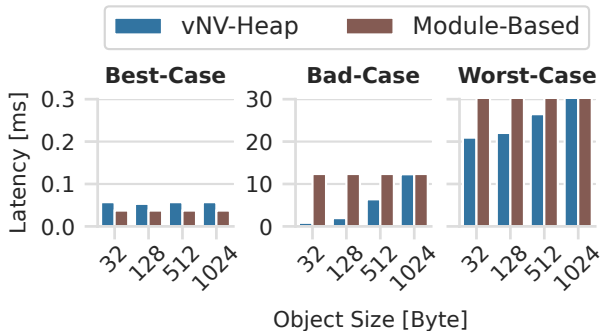
Evaluation Setup



Reference Retrieval



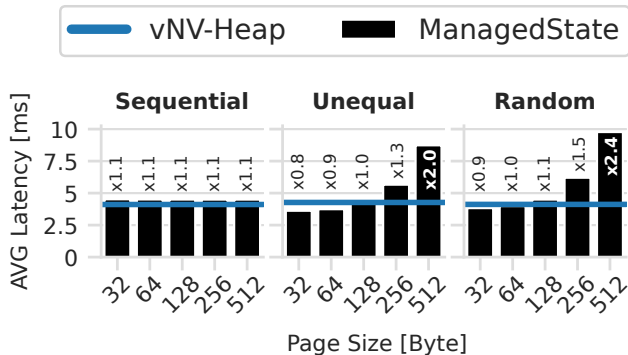
Reference Retrieval



Takeaway #2

More **fine grained management** of objects **reduces swapping overhead up to 93%** in the worst/relevant case!

Key-Value Store



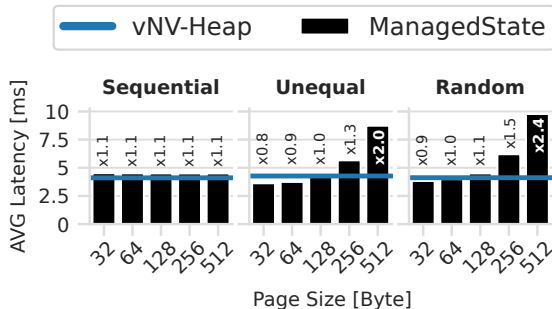
KVS Keys

32-bit integers

KVS Values (58 KiB)

- 64×32 bytes
- 128×128 bytes
- 32×256 bytes
- 32×1024 bytes

Key-Value Store

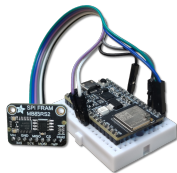


Takeaway #3

The vNV-Heap provides **similar/balanced performance**, while providing more features: **NVM safety & transparent swapping!**

Conclusion

vNV-Heap: virtually Non-Volatile Hheap



NVM safety



Transparent swapping

→ Treat memory as resource & guard it with ownership/borrowing



Predictable

→ Provides worst-case guarantees

→ Improves the worst case with object-wise modification tracking



HW agnostic

→ Implemented & employable as a software library

Thank you for listening!



Link to Paper



Link to Code

References (1)

- [1] M. Afanasov et al. **“Battery-less zero-maintenance embedded sensing at the mithræum of circus maximus”**. In: SenSys '20. Virtual Event, Japan: Association for Computing Machinery, 2020, pp. 368–381. ISBN: 9781450375900. DOI: [10.1145/3384419.3430722](https://doi.org/10.1145/3384419.3430722).
- [2] S. Sliper et al. **“Efficient State Retention through Paged Memory Management for Reactive Transient Computing”**. In: DAC '19. ACM, June 2019. DOI: [10.1145/3316781.3317812](https://doi.org/10.1145/3316781.3317812).

- [3] J. de Winkel, H. Tang, and P. Pawełczak. **“Intermittently-powered bluetooth that works”**. In: MobiSys '22. Portland, Oregon, 2022, pp. 287–301. DOI: [10.1145/3498361.3538934](https://doi.org/10.1145/3498361.3538934).